List Distinguishing Parameters of Trees

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Abstract

A coloring of the vertices of a graph G is said to be *distinguishing* provided no nontrivial automorphism of G preserves all of the vertex colors. The *distinguishing number* of G, D(G), is the minimum number of colors in a distinguishing coloring of G. The *distinguishing chromatic number* of G, $\chi_D(G)$, is the minimum number of colors in a distinguishing coloring of G that is also a proper coloring.

Recently the notion of a distinguishing coloring was extended to that of a list distinguishing coloring. Given an assignment $L = \{L(v)\}_{v \in V(G)}$ of lists of available colors to the vertices of G, we say that G is (properly) *L*-distinguishable if there is a (proper) distinguishing coloring f of G such that $f(v) \in L(v)$ for all v. The list distinguishing number of G, $D_{\ell}(G)$, is the minimum integer k such that G is *L*-distinguishable for any list assignment L with |L(v)| = k for all v. Similarly, the list distinguishing chromatic number of G, denoted $\chi_{D_{\ell}}(G)$ is the minimum integer k such that G is properly *L*-distinguishable for any list assignment L with |L(v)| = k for all v.

In this paper, we study these distinguishing parameters for trees, and in particular extend an enumerative technique of Cheng to show that for any tree T, $D_{\ell}(T) = D(T)$, $\chi_D(T) = \chi_{D_{\ell}}(T)$, and $\chi_D(T) \leq D(T) + 1$.

Keywords: Distinguishing Coloring, List Distinguishing Coloring, Proper Distinguishing Coloring, Distinguishing Chromatic Number, List Distinguishing Chromatic Number

1 Introduction

A coloring of a graph G is a labeling $\phi : V(G) \to \mathbb{N}$; a k-coloring is a labeling $\phi : V(G) \to [k]$, where $[k] = \{1, 2, \dots, k\}$. A coloring of the vertices of a graph G is distinguishing if no nontrivial automorphism of G preserves all of the vertex colors; such a coloring distinguishes G. In 1996 Albertson and

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Collins [1] introduced the distinguishing number D(G) of a graph G, defined to be the minimum number of colors in a distinguishing coloring of G. Since its introduction, the distinguishing number and related parameters have received considerable attention in the literature (see for example [4, 5, 12, 15, 16]). In 2006 Collins and Trenk [9] introduced the distinguishing chromatic number $\chi_D(G)$ of a graph G, defined to be the minimum number of colors in a distinguishing coloring of G that is also a proper coloring. Subsequent investigation of the distinguishing chromatic number (for instance, [3, 8, 13, 17]) has focused on similarities and disparities between the distinguishing chromatic number and one or both of the distinguishing number and the chromatic number.

One of the most studied variants of the chromatic number is the list chromatic number, introduced independently by Vizing in 1976 [18] and Erdős, Rubin, and Taylor in 1979 [10]. Recently Ferrara, Flesch, and Gethner [11] extended the notion of a distinguishing coloring to that of a list distinguishing coloring. Given an assignment $L = \{L(v)\}_{v \in V(G)}$ of lists of available colors to the vertices of G, we say that G is L-distinguishable if there is a distinguishing coloring f of G such that $f(v) \in L(v)$ for all v. The list distinguishing number of G, $D_{\ell}(G)$, is the minimum integer ksuch that G is L-distinguishable for any list assignment L with |L(v)| = k for all v. The existence of a graph G for which D(G) and $D_{\ell}(G)$ are not equal remains a major open question.

While not explicitly introduced in [11], it is natural to consider a list analogue of the distinguishing chromatic number. We say that G is properly L-distinguishable if there is a distinguishing coloring f of G chosen from the lists such that f is also a proper coloring of G. The list distinguishing chromatic number $\chi_{D_{\ell}}(G)$ of G is the minimum integer k such that G is properly L-distinguishable for any assignment L of lists with $|L(v)| \geq k$ for all v.

In this paper, we study these four distinguishing parameters for trees, showing that $D_{\ell}(T) = D(T)$, $\chi_D(T) = \chi_{D_{\ell}}(T)$, and $\chi_D(T) \leq D(T) + 1$ for any tree T. In our proofs, we extend an enumerative technique first introduced by Cheng to determine the distinguishing number of trees [6]. The technique has subsequently been used to determine the distinguishing number of both planar graphs [2] and interval graphs [7].

1.1 Preliminaries

A vertex v is in the *center* of a tree T if v minimizes $\max_{u \in V(T)} \{ \text{dist}(u, v) \}$. In the majority of the paper, we will be considering rooted trees. A *rooted tree* is a tree with an identified vertex called the *root*. Given vertices u and v in a rooted tree T with root r, v is the *parent* of u if v is the neighbor of u in the unique u, r-path in T, and u in this case is a *child* of v. Two vertices with the same parent are *siblings*. More generally, v is an *ancestor* of u if v lies anywhere in the unique u, r-path, and in this case u is a *descendant* of v. Given a vertex v in a rooted tree T, we let T_v denote the subtree of T rooted at v that is induced by v and all of its descendants. All automorphisms of a rooted tree stabilize the root.

Two colorings ϕ_1 and ϕ_2 of rooted trees T_{v_1} and T_{v_2} are *equivalent* if there exists an isomorphism $\pi : V(T_{v_1}) \to V(T_{v_2})$ that maps v_1 to v_2 and that maps the coloring ϕ_1 to ϕ_2 : that is, for every vertex $u \in V(T_{v_1})$, we have $\phi_1(u) = \phi_2(\pi(u))$.

2 Distinguishing Colorings

The main result of this section is the following.

Theorem 1. If T is a tree, then $D(T) = D_{\ell}(T)$.

To prove Theorem 1, we shall prove $D(T') = D_{\ell}(T')$ for any rooted tree T'; this suffices since the following lemma provides a reduction from unrooted trees to rooted trees.

Lemma 2. For any tree T, there is a rooted tree T' with D(T) = D(T') and $D_{\ell}(T) = D_{\ell}(T')$.

Proof. In [14], Jordan proved that the center of a tree T is either a single vertex or an edge, and in either case the center of T is set-wise stabilized by all automorphisms. If the center of T is a single vertex, then let T' be T rooted at the center and note that Aut(T) = Aut(T').

If the center of T is the edge e, then let T' be obtained by subdividing e and rooting the resulting tree at the unique vertex in T' - T, which is the center of T'. In [6], Cheng showed that D(T) = D(T'), and it is straightforward to show by a nearly identical argument that $D_{\ell}(T) = D_{\ell}(T')$ as well.

The following observation describes a simple condition that is necessary and sufficient for a coloring to distinguish a rooted tree.

Observation 3. Let T be a rooted tree. For a fixed $x \in V(T)$, partition the children of x into classes C_1, \ldots, C_t so that two vertices u, v are in the same class if and only if $T_u \cong T_v$. A coloring ϕ distinguishes T_x if and only if

- 1. for all children v of x, the restriction of ϕ to T_v distinguishes T_v , and
- 2. given vertices $u, v \in C_j$ for $j \in [t]$, the restrictions of ϕ to T_u and T_v are inequivalent.

Following the notation from [6], we let D(G;k) denote the number of equivalence classes of distinguishing k-colorings of a graph G. Cheng showed that $D(T_x;k)$ can be computed recursively.

Lemma 4 (Cheng [6]). Let T_x be a tree with root x and partition the children of x into classes C_1, \ldots, C_t so that two vertices u, v are in C_j if and only if $T_u \cong T_v$. Select a representative u_j for each class C_j . The number of equivalence classes of distinguishing k-colorings of T_x is given by

$$D(T_x;k) = k \prod_{j=1}^t \binom{D(T_{u_j};k)}{|C_j|}.$$

For a list assignment $L = \{L(v)\}_{v \in V(G)}$, we let D(G; L) be the number of equivalence classes of distinguishing L-colorings of G.

Theorem 5. If T_x is a rooted tree and $L = \{L(v)\}_{v \in V(T_x)}$ is a list assignment with |L(v)| = k for all $v \in V(T_x)$, then $D(T_x; L) \ge D(T_x; k)$. Equality holds if and only if $D(T_x; L) = 0$ or L(u) = L(v) whenever u and v lie in the same orbit of T_x .

Proof. Observe first that if $D(T_x; k) = 0$, then the conclusion holds trivially. For any vertex $v \in T$, let L_v denote the restriction of L to the vertices in T_v . We proceed by induction to show that $D(T_v; L_v) \ge D(T_v; k)$ for each subtree T_v of T_x . If v is a leaf, then $D(T_v, L_v) = D(T_v; k) = k$.

Assume that v is not a leaf of T_x . Partition the children of v into classes C_1, \ldots, C_t as in Observation 3, and label the vertices in class C_j as $v_{j,1}, v_{j,2}, \ldots, v_{j,m_j}$ where $m_j = |C_j|$. By induction, $D(T_{v_{j,p}}; L_{v_{j,p}}) \ge D(T_{v_{j,p}}; k)$ for all $p \in [m_j]$. Set $d_j = D(T_{v_{j,1}}; k)$; it follows that $d_j = D(T_{v_{j,p}}; k)$ for all $p \in [m_j]$ since the subtrees in $T_{v_{j,p}}$ are isomorphic.

For $p \in [m_j]$, let $S_{j,p}$ be a set of representative colorings from each equivalence class of distinguishing $L_{v_{j,p}}$ -colorings of $T_{v_{j,p}}$; note that $|S_{j,p}| = D(T_{v_{j,p}}; L_{v_{j,p}})$, and by induction, $|S_{j,p}| \ge d_j$. Let R_j be the set of tuples $(\phi_{j,1}, \phi_{j,2}, \ldots, \phi_{j,m_j}) \in S_{j,1} \times S_{j,2} \times \cdots \times S_{j,m_j}$ such that $\phi_{j,p}$ and $\phi_{j,q}$ are inequivalent when $p \neq q$. By Observation 3, an L_v -coloring ϕ distinguishes T_v if and only if, for every class C_j , the tuple $(\phi_{j,1}, \phi_{j,2}, \ldots, \phi_{j,m_j})$ of colorings induced on $T_{v_{j,p}}$ is component-wise equivalent to an element of R_j . Let r_j be the number of equivalence classes in R_j , where two tuples are equivalent if they are the same up to the order of the coordinates. A maximum set of inequivalent L_v -colorings of T_v is formed by independently selecting one tuple $(\phi_{j,1}, \phi_{j,2}, \ldots, \phi_{j,m_j})$ from each equivalence class of R_j for each class C_j and selecting any color for v from L(v). Hence, $D(T_v; L_v) = k \prod_{j=1}^t r_j$.

We bound r_j by selecting colorings $(\phi_{j,1}, \phi_{j,2}, \ldots, \phi_{j,m_j})$ for the subtrees $T_{v_{j,1}}, \ldots, T_{v_{j,m_j}}$ in order. By induction, for $p \in [m_j]$, there are at least $d_j - p + 1$ choices for $\phi_{j,p}$ in $S_{j,p}$ that are inequivalent to the previous selections for $\phi_{j,1}, \ldots, \phi_{j,p-1}$. Thus, there are at least $d_j(d_j - 1) \cdots (d_j - m_j + 1)$ ways to select m_j inequivalent colorings $(\phi_{j,1}, \phi_{j,2}, \ldots, \phi_{j,m_j})$. Each equivalence class in R_j is counted at most m_j ! times, so

$$r_j \ge \frac{d_j(d_j-1)\cdots(d_j-m_j+1)}{m_j!} = \binom{d_j}{m_j} = \binom{D(T_{v_{j,1}};k)}{|C_j|}.$$

Therefore by Lemma 4,

$$D(T_v; L_v) = k \prod_{j=1}^t r_j \ge k \prod_{j=1}^t {D(T_{v_{j,1}}; k) \choose |C_j|} = D(T_v; k).$$

If $D(T_x; L) = 0$, then it is clear that $D(T_x; L) = D(T_x; k)$. Otherwise we show that equality holds if and only if L(u) = L(v) when u and v lie in the same orbit of T_x . We prove this by induction on the number of vertices in T_x . Let L be a list assignment such that $D(T_x; L) > 0$ and $D(T_x; L) = D(T_x; k)$. The result holds trivially if T_x has a single vertex.

Let C_1, \ldots, C_t be the partition of the children of x as above. Equality holds for T_x if and only if $r_j = \binom{D(T_{v_{j,1}};k)}{|C_j|}$ for all $j \in [t]$. Furthermore, $r_j = \binom{D(T_{v_{j,1}};k)}{|C_j|} = \frac{d_j(d_j-1)\cdots(d_j-m_j+1)}{m_j!}$ if and only if $|S_{j,p}| = d_j$ for all $p \in [m_j]$ and every coloring in $S_{j,p}$ has an equivalent coloring in $S_{j,q}$ for all $p, q \in [m_j]$. By induction, $|S_{j,p}| = d_j$ if and only if L(y) = L(w) when y and w lie in the same orbit of $T_{v_{j,p}}$.

Consider a vertex y in $T_{v_{j,p}}$. By permuting colors on the orbit 0 containing y, each color in L(y)appears in 0 in a distinguishing L-coloring $\phi_{j,p}$ in $S_{j,p}$. The isomorphism $\sigma : V(T_{v_{j,p}}) \to V(T_{v_{j,q}})$ guaranteed by an equivalent pair $(T_{v_{j,p}}, \phi_{j,p}) \cong (T_{v_{j,q}}, \phi_{j,q})$ satisfies $\phi_{j,p}(y) = \phi_{j,q}(\sigma(y))$ for all $y \in V(T_{v_{j,p}})$. Therefore, each color in L(y) also appears in the lists of the vertices in the orbit of $\sigma(y)$ in $T_{v_{j,q}}$, and consequently $L(y) = L(\sigma(y))$. Each vertex in the orbit of y in T_x that is not in $T_{v_{j,p}}$ lies in the image of \mathcal{O} under some such isomorphism σ . Thus, every vertex in the orbit of y in T_x has the same list of colors as y. Consequently, if equality holds, then L(u) = L(v) whenever u and v lie in the same orbit of T_x .

Conversely, if L(u) = L(v) whenever u and v lie in the same orbit of T_x , then L(y) = L(w) when y and w lie in the same orbit of $T_{v_{j,p}}$. Furthermore, each coloring $S_{j,p}$ has an equivalent coloring in $S_{j,q}$ for all $p, q \in [m_j]$.

Note that D(T) is the minimum k so that D(T; k) is positive. Similarly, $D_{\ell}(T)$ is the minimum k so that D(T; L) is positive for every list assignment L with |L(v)| = k. By considering the list assignment that gives the same k colors to every vertex, it is clear that $D_{\ell}(T)$ is positive only if D(T) is. Thus, the following corollary is immediate from Theorem 5.

Corollary 6. If T is a rooted tree, then $D(T) = D_{\ell}(T)$.

Theorem 1 follows from Lemma 2 and Corollary 6.

3 List Distinguishing Chromatic Number

In this section, we prove that the distinguishing chromatic number of a tree is equal to the list distinguishing chromatic number using a similar enumerative method.

Theorem 7. If T is a tree, then $\chi_D(T) = \chi_{D_\ell}(T)$.

For $x \in V(T)$, let $D_{\chi}(T_x; k)$ denote the number of equivalence classes of distinguishing proper k-colorings of the rooted tree T_x . For $i \in [k]$, let $D_{\chi}(T_x; k, i)$ denote the number of equivalence classes of distinguishing proper k-colorings of T_x in which x gets color i. Note that $D_{\chi}(T_x; k, i) = D_{\chi}(T_x; k, 1) = \frac{1}{k} D_{\chi}(T_x; k)$ for all $i \in [k]$. Hence $D_{\chi}(T_x; k) = k D_{\chi}(T_x; k, 1)$.

Similarly, let $D_{\chi}(T_x; L)$ denote the number of equivalence classes of distinguishing proper *L*-colorings of the rooted tree T_x . For $i \in L(x)$, let $D_{\chi}(T_x; L, i)$ be the number of equivalence classes of distinguishing proper *L*-colorings of T_x in which x gets color i. Here, the value $D_{\chi}(T_x; L, i)$ may change for different values of $i \in L(x)$. Thus $D_{\chi}(T_x; L) = \sum_{i \in L(x)} D_{\chi}(T_x; L, i)$.

As we discuss below, consideration of rooted trees is sufficient to demonstrate $\chi_D(T) = \chi_{D_\ell}(T)$ in most, but not all unrooted cases. The next result is analogous to Theorem 5.

Theorem 8. If T_x is a rooted tree and $L = \{L(v)\}_{v \in V(T_x)}$ is a list assignment with |L(v)| = k for all $v \in V(T_x)$, then $D_{\chi}(T_x; L, i) \ge D_{\chi}(T_x; k, 1)$ for all $i \in L(x)$. Equality holds for all $i \in L(x)$ if and only if $k \ge 2$ and all lists in L are identical.

Proof. Similar to the proof of Theorem 5, we show that $D_{\chi}(T_v; L_v, i) \ge D_{\chi}(T_v; k, 1)$ for all $v \in V(T_x)$ and $i \in L(v)$ inductively. If v is a leaf, then $D_{\chi}(T_v; L_v, i) = D_{\chi}(T_v; k, 1) = 1$.

Assume that v is not a leaf. Partition the children of v into equivalence classes C_1, C_2, \ldots, C_t by isomorphism classes of subtrees. Label the vertices in each class C_j as $v_{j,1}, v_{j,2}, \ldots, v_{j,m_j}$ where $m_j = |C_j|$. The following claim is analogous to Lemma 4.

Claim 8.1. For
$$i \in [k]$$
, $D_{\chi}(T_v; k, i) = \prod_{j=1}^{t} \binom{(k-1)D_{\chi}(T_{v_{j,1}}; k, 1)}{|C_j|}.$

Proof. By Observation 3, a proper k-coloring ϕ such that $\phi(v) = i$ distinguishes T_v if and only if the colorings $\phi_{v_{j,p}}$, obtained by restricting ϕ to the trees $T_{v_{j,p}}$ for $v_{j,p} \in C_j$, are inequivalent distinguishing proper k-colorings with $\phi_{v_{j,p}}(v_{j,p}) \neq i$. Hence, there are $(k-1)D_{\chi}(T_{v_{j,1}};k,1)$ possible colorings for $T_{v_{j,p}}$. Since $T_{v_{j,p}} \cong T_{v_{j,1}}$ for all $p \in [m_j]$, there are $(k-1)D_{\chi}(T_{v_{j,1}};k,1)$ possible colorings of $T_{v_{j,p}}$. We choose $|C_j|$ of these colorings to place on the trees $T_{v_{j,p}}$ for $v_{j,p} \in C_j$. Since $D_{\chi}(T_v;k,i)$ counts colorings up to isomorphism, the selection of these colorings is independent of their order in C_j . Also, the choices of these colorings are independent among the different isomorphism classes $\{C_j\}_{j=1}^t$, so $D_{\chi}(T_v;k,i)$ is given by the product $\prod_{j=1}^t {\binom{(k-1)D_{\chi}(T_{v_{j,1}};k,1)}{|C_j|}}$, proving the claim. \Box

Let $d_{j}^{i} = (k-1)D_{\chi}(T_{v_{j,1}}; k, 1)$. By induction,

$$\sum_{i' \in L(v_{j,p}) \setminus \{i\}} D_{\chi}(T_{v_{j,p}}; L_{v_{j,p}}, i') \ge (k-1)D_{\chi}(T_{v_{j,p}}; k, 1) = d_j^i$$

We now consider the number of equivalence classes of distinguishing proper $L_{v_{j,p}}$ -colorings of $T_{v_{j,p}}$ in which $v_{j,p}$ does not get color *i*. Let $S_{j,p}^i$ be a set of representative colorings from each such equivalence class. Let R_j^i be the set of tuples $(\phi_{j,1}, \phi_{j,2}, \ldots, \phi_{j,m_j}) \in S_{j,1}^i \times S_{j,2}^i \times \cdots \times S_{j,m_j}^i$ such that $\phi_{j,p}$ and $\phi_{j,q}$ are inequivalent when $p \neq q$. By induction,

$$|S_{j,p}^{i}| = \sum_{i' \in L(v_{j,p}) \setminus \{i\}} D_{\chi}(T_{v_{j,p}}; L_{v_{j,p}}, i') \ge d_{j}^{i}.$$

By Observation 3, a proper L_v -coloring distinguishes T_v if and only if, for every class C_j , the colorings $\phi_{j,p}$ induced on $T_{v_{j,p}}$ form a tuple $(\phi_{j,1}, \phi_{j,2}, \ldots, \phi_{j,m_j}) \in R_j^i$. Let r_j^i be the number of equivalence classes in R_j^i , where two tuples are equivalent if they are the same up to the order of the coordinates. Form a maximum set of inequivalent proper L_v -colorings ϕ in which v gets color i by independently selecting one tuple $(\phi_{j,1}, \phi_{j,2}, \ldots, \phi_{j,m_j})$ from each equivalence class of R_j^i for each class C_j . Hence, $D_{\chi}(T_v; L_v, i) = \prod_{j=1}^t r_j^i$.

We bound r_j^i by selecting colorings $(\phi_{j,1}, \ldots, \phi_{j,m_j})$ for subtrees in order. For $p \in [m_j]$, there are at least $d_j^i - p + 1$ selections for $\phi_{j,p}$ from $S_{j,p}$ that are inequivalent to the selections of $\phi_{j,1}, \ldots, \phi_{j,p-1}$. Thus, there are at least $d_j^i(d_j^i - 1) \cdots (d_j^i - m_j + 1)$ possible ways to select m_j inequivalent colorings $\phi_{j,1}, \phi_{j,2}, \ldots, \phi_{j,m_j}$. Each equivalence class in R_j^i is counted at most m_j ! times, so

$$r_j^i \ge \frac{d_j^i(d_j^i - 1) \cdots (d_j^i - m_j + 1)}{m_j!} = \binom{d_j^i}{m_j} = \binom{(k-1)D_{\chi}(T_{v_{j,1}}; k, 1)}{|C_j|}.$$

Therefore, by Claim 8.1,

$$D_{\chi}(T_v; L, i) = \prod_{j=1}^{t} r_j^i \ge \prod_{j=1}^{t} \binom{(k-1)D_{\chi}(T_{v_{j,1}}; k, 1)}{|C_j|} = D_{\chi}(T_v; k, 1).$$

Consider the partition of the children of x as in Observation 3. Equality holds for all colors in L(x) if and only if, given an arbitrary $i \in L(x)$, $r_j^i = \binom{(k-1)D_{\chi}(T_{v_{j,1}};k,1)}{|C_j|}$ for all for all $j \in [t]$; this holds if and only if $|S_{j,p}^i| = d_j^i = (k-1)D_{\chi}(T_{v_{j,1}};k;1)$ for all $p \in [m_j]$. Noting that $(k-1)D_{\chi}(T_{v_{j,1}};k;1) = (k-1)D_{\chi}(T_{v_{j,p}};k;1)$, equality holds if and only if $D_{\chi}(T_{v_{j,p}};k;1) = D_{\chi}(T_{v_{j,p}};L;i')$ for all $i' \in L(v_{j,p}) \setminus \{i\}$. Because there are multiple choices of $i \in L(x)$, it follows that equality holds for all colors in L(x) if and only if $D_{\chi}(T_{v_{j,p}};k;1) = D_{\chi}(T_{v_{j,p}};L;i')$ for all $i' \in L(v_{j,p})$. Thus, by induction, the lists are identical in $T_{v_{j,p}}$. Furthermore, $r_j^i = \binom{(k-1)D_{\chi}(T_{v_{j,1}};k,1)}{|C_j|}$ for all $j \in [t]$ if and only if color i appears in the list $L(v_{j,p})$ for all $p \in [m_j]$. This is true for all $i \in L(x)$, so x and all of its children have the same list.

If T_x is a rooted tree, then $\chi_D(T_x)$ is the minimum k so that $D_{\chi}(T_x; k, 1) > 0$, and $\chi_{D_{\ell}}(T_x)$ is the minimum k so that for every assignment L of lists of size k, there exists an $i \in L(x)$ such that $D_{\chi}(T_x; L, i)$ is positive. By considering the list assignment in which every vertex gets the list [k], the following corollary is immediate from Theorem 8.

Corollary 9. If T_r is a rooted tree, then $\chi_D(T_r) = \chi_{D_\ell}(T_r)$.

The proof of Theorem 1 followed from the reduction from unrooted to rooted trees in Lemma 2 and the equality given in Corollary 6. For the distinguishing chromatic number and list distinguishing chromatic number, such a reduction works in most cases.

Lemma 10. If the center of T is a single vertex or $\chi_D(T) \ge 3$, then $\chi_D(T) = \chi_{D_\ell}(T)$.

Proof. Case 1: The center of T is a vertex x. Since all automorphisms set-wise stabilize the center, x is stabilized by all automorphisms of T. Rooting the tree at x does not change the automorphism group, so a T_x -distinguishing proper coloring of V(T) is also a T-distinguishing proper coloring. If T' is T rooted at x, then $\chi_D(T) = \chi_D(T') = \chi_{D_\ell}(T') = \chi_{D_\ell}(T)$.

Case 2: The center of T is an edge uv. Let T' be obtained by subdividing uv and rooting the resulting tree at the unique vertex x in T' - T, which is the center of T'. The automorphisms of T' are given by the actions of automorphisms of T on V(T) while stabilizing x.

By Corollary 9, $\chi_D(T') = \chi_{D_\ell}(T')$. Since any *T*-distinguishing proper *k*-coloring ϕ uses at least 3 colors, it is possible to extend ϕ to a *T'*-distinguishing proper *k*-coloring by assigning *x* a color that is not $\phi(u)$ or $\phi(v)$. Thus $\chi_D(T') \leq \chi_D(T)$.

Since T' contains an edge, $\chi_{D_{\ell}}(T') \geq 2$; let $k = \chi_{D_{\ell}}(T')$. Consider a list assignment L on V(T')in which every list has size k, and let L_T be the restriction of L to V(T). Let T'_v be the subtree of T' rooted at v consisting of all descendants of v; clearly $\chi_{D_{\ell}}(T'_v) \leq \chi_{D_{\ell}}(T')$. By Corollary 9, $\chi_D(T'_v) = \chi_{D_{\ell}}(T'_v)$, so there is a T'_v -distinguishing proper k-coloring and we may assume that vreceives color 1 in such a coloring. Thus, by Theorem 8, there is a T'_v -distinguishing proper L_v coloring in which v gets color i for each $i \in L(v)$. Similarly, there is a T'_u -distinguishing proper L_u coloring in which u gets color j for each $j \in L(u)$. Because all lists have size at least 2, it is possible to choose colors $i \in L(v)$ and $j \in L(u)$ so that $i \neq j$. The combination of a distinguishing proper L_v -coloring of T'_v in which v receives color i and a distinguishing proper L_u -coloring of T'_u in which u receives color j produces a distinguishing proper L_T -coloring of T. Since every list assignment on V(T) is the restriction of some list assignment on V(T'), it follows that $\chi_{D_\ell}(T) \leq \chi_{D_\ell}(T')$. By considering the list assignment in which every vertex gets the same list of length k, we see that $\chi_D(T) \leq \chi_{D_\ell}(T)$. Therefore

$$\chi_D(T) \le \chi_{D_\ell}(T) \le \chi_{D_\ell}(T') = \chi_D(T') \le \chi_D(T).$$

Lemma 11. If T is an unrooted tree with center $\{u, v\}$ and $\chi_D(T) = 2$, then $\chi_{D_\ell}(T) = 2$.

Proof. Let ϕ be a proper 2-coloring of T. Note that $\phi(u) \neq \phi(v)$. Since T is connected, the color at any vertex $y \in V(T)$ is specified by the parity of the distance from y to u: $\phi(y) = \phi(u)$ if and only if the distance from y to u is even. Hence, there are exactly two proper 2-colorings of T and they are equivalent under a permutation of the colors. Since $\chi_D(T) = 2$, such a coloring must distinguish T. The center of T is set-wise stabilized by all automorphisms of T, so the distance from any vertex $y \in V(T)$ to $\{u, v\}$ is preserved under all automorphisms of T. If u and v are point-wise stabilized by a non-trivial automorphism σ , then every $y \in V(T)$ is mapped to another element of the same distance to u. Therefore $\phi(y) = \phi(\sigma(y))$ and this 2-coloring does not distinguish T. Thus, any non-trivial automorphism ϕ must swap u and v. The product of any two non-trivial automorphisms point-wise stabilize u and v, and therefore the product is the identity. Hence, the automorphism group of T is isomorphic to \mathbb{Z}_2 and any coloring ϕ where $\phi(u) \neq \phi(v)$ distinguishes T. Thus, $\chi_{D_\ell}(T) = 2 = \chi_D(T)$.

Theorem 7 follows immediately from Lemma 10 and Lemma 11.

4 Distinguishing Chromatic Number

In this section, we show that the distinguishing chromatic number and the distinguishing number of a tree differ by at most one and the bound is sharp. We also describe when these parameters are different.

Theorem 12. If T_x is a rooted tree, then $\chi_D(T_x) \leq D(T_x) + 1$.

Proof. Let $\phi: V(T_x) \to \{1, \ldots, k\}$ be a distinguishing k-coloring of T_x . We create a distinguishing proper (k + 1)-coloring $\phi': V(T) \to \{1, \ldots, k, *\}$ starting at the root: let $\phi'(x) = \phi(x)$. Proceed recursively; after coloring a vertex v, consider a child u of v. If $\phi'(v) = \phi(u)$, then let $\phi'(u) = *$. Otherwise, let $\phi'(u) = \phi(u)$. Hence, a vertex u receives color * if and only if its parent v has $\phi(v) = \phi(u)$ and the previous step assigned $\phi'(v) = \phi(v)$. Therefore, ϕ' is a proper coloring. Also note that if x and y are siblings with $\phi'(x) = \phi'(y)$, then $\phi(x) = \phi(y)$. Since any nontrivial automorphism of T_x interchanges subtrees rooted at siblings, any ϕ' -preserving automorphism is also a ϕ -preserving automorphism. Therefore ϕ' is distinguishing if ϕ is distinguishing. \Box

It follows from Lemmas 2, 10, and 11 that $\chi_D(T) \leq D(T) + 1$ for unrooted trees as well.

We now characterize the trees for which $\chi_D(T) = D(T) + 1$. Again we let T' be the rooted tree obtained by rooting T at its center if the center is unique, or rooting at the vertex obtained by subdividing the central edge of T.

Theorem 13. If T is a tree with D(T) = k, then $\chi_D(T) = k + 1$ if and only if

- 1. $|V(T)| \ge 2$ and k = 1, or
- 2. there exists a vertex $x \in V(T)$ and a set S of children of x in T' so that the subtrees T_u and T_v are isomorphic for all $u, v \in S$ and $(k-1)D_{\chi}(T_u; k, 1) < |S|$ for $u \in S$.

Proof. If D(T) = 1 and T has more than one vertex, then clearly $\chi_D(T) = 2$. Thus we may assume that $D(T) \ge 2$, and we only consider graphs satisfying $\chi_D(T) \ge 3$. By Lemmas 2 and 10, it suffices to consider the rooted tree T'.

By Claim 8.1, the number of proper distinguishing k-colorings of the rooted tree T'_x is given by $kD_{\chi}(T'_x;k,1) = k\prod_{j=1}^t {\binom{(k-1)D_{\chi}(T'_{v_j};k,1)}{|C_j|}}$. There exists no proper distinguishing k-coloring of T' if and only if this product is zero. This product is zero if and only if $(k-1)D_{\chi}(T'_{v_j};k,1) < |C_j|$ for some j. If $(k-1)D_{\chi}(T'_{v_j};k,1) < |C_j|$ for some j, then letting $S = C_j$ suffices. Conversely, if the set S exists, then $S \subseteq C_j$ for some j. If v_j is an element of S, then $(k-1)D_{\chi}(T'_{v_j};k,1) < |C_j|$, and there is no proper distinguishing k-coloring of T'.

We conclude this section by noting that Cheng's algorithm for computing D(T) can be adapted to compute $\chi_D(T)$ by replacing the counting method in Lemma 4 with Claim 8.1. This leads to a polynomial time¹ algorithm to determine if $\chi_D(T) = D(T)$ and if not, it can produce the certificate (v, S) from Theorem 13.

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¹Cheng's algorithm runs in $O(n \log n)$ time and this modification replaces the formula in Lemma 4 with the formula in Claim 8.1, giving $O(n \log n)$ time to compute $\chi_D(T)$.

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