

Rainbow Arithmetic Progressions

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`http://www.math.iastate.edu/dstolee/r/rainbowaps.htm`

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Ramsey Theory and anti-Ramsey Theory

Ramsey Theory: *Looking for monochromatic (mono_χ) subgraphs in large edge-colored graphs.*

Anti-Ramsey Theory: *Looking for rainbow subgraphs in edge-colorings using many colors.*

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“Complete disorder is unavoidable.”

Ramsey Theory on the Integers

We will consider $[n] = \{1, \dots, n\} \subset \mathbb{N}$. Let $k \geq 3$.

Definition A **k -term arithmetic progression (k -AP)** is a set S such that

$$S = \{a + id : 0 \leq i < k\} = \{a, a + d, a + 2d, \dots, a + (k - 1)d\}$$

for some integers a, d , and $d \neq 0$.

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van der Waerden: *If the number of colors r is fixed and n is large, then there exists a monochromatic k -AP in every $c : [n] \rightarrow [r]$.*

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van der Waerden: *If the number of colors r is fixed and n is large, then there exists a mono_χ k -AP in every $c : [n] \rightarrow [r]$.*

Definition $w_r(k)$ is the minimum n such that all r -colorings of $[n]$ contain a mono_χ k -AP.

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$$\lim_{n \rightarrow \infty} \text{sz}([n], k) / n = 0.$$

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Theorem (Behrend)

$$\text{sz}([n], k) \geq ne^{-b\sqrt{\log n}}.$$

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Theorem (Juncić, Fox, Mahdian, Nešetřil, Radoičić) If \mathbb{N} is colored with three colors such that each color class has upper density strictly larger than $\frac{1}{6}$, then the coloring contains a rainbow 3-AP.

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Theorem (Axenovich, Fon-Der-Flaass) If $[n]$ is colored with three colors such that each color class has size at least $\frac{n+4}{6}$, then the coloring contains a rainbow 3-AP.

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An r -coloring is **exact** if all colors are used at least once.

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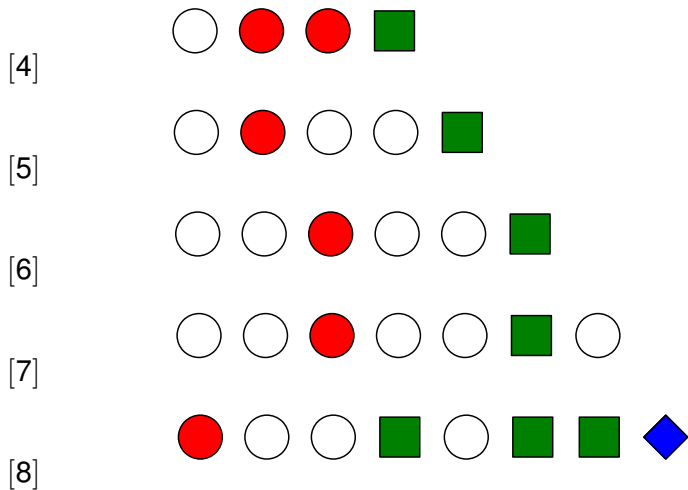
Assuming $k \leq n$:

$$k \leq aw([n], k) \leq n$$

$n \backslash k$	3	4	5	6	7	8	9
3	3						
4	4						
5	4	5					
6	4	6					
7	4	6	7				
8	5	6	8				
9	4	7	8	9			
10	5	8	9	10			
11	5	8	9	10	11		
12	5	8	10	11	12		
13	5	8	11	11	12	13	
14	5	8	11	12	13	14	
15	5	9	11	13	14	14	15

Values of $\text{aw}([n], k)$ for $3 \leq k \leq \frac{n+3}{2}$.

Extremal Colorings with no Rainbow 3-AP



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[4]



[5]



[6]



[7]



[8]



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[4]



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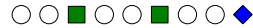
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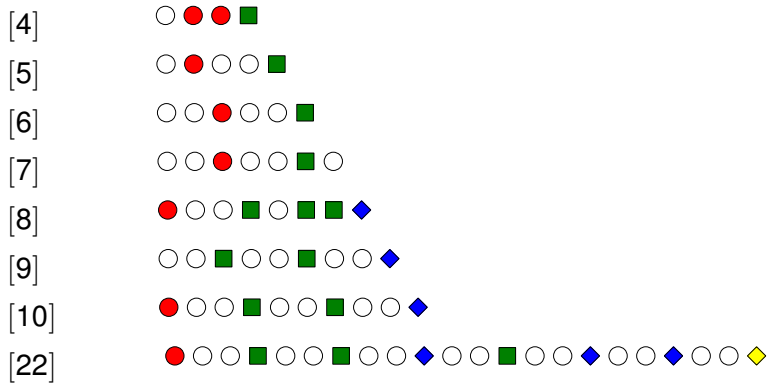
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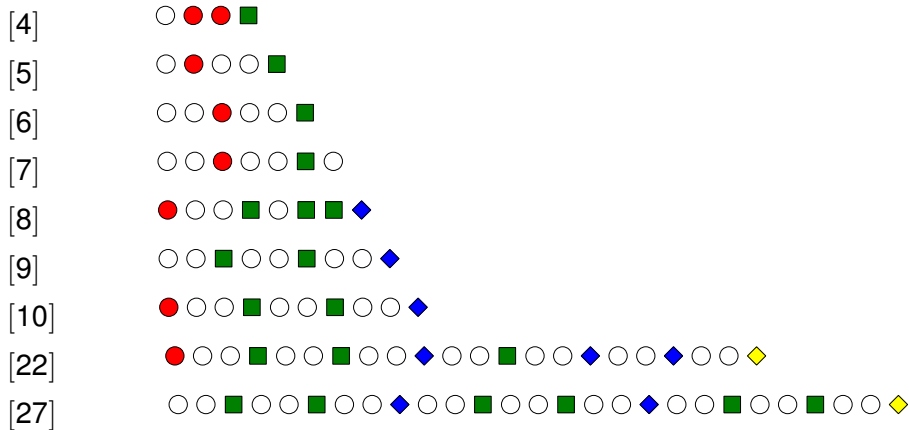
Extremal Colorings with no Rainbow 3-AP

[4]	○ ● ● ■
[5]	○ ● ○ ○ ■
[6]	○ ○ ● ○ ○ ■
[7]	○ ○ ● ○ ○ ■ ○
[8]	● ○ ○ ■ ○ ■ ■ ◆
[9]	○ ○ ■ ○ ○ ■ ○ ○ ◆
[10]	● ○ ○ ■ ○ ○ ■ ○ ○ ◆

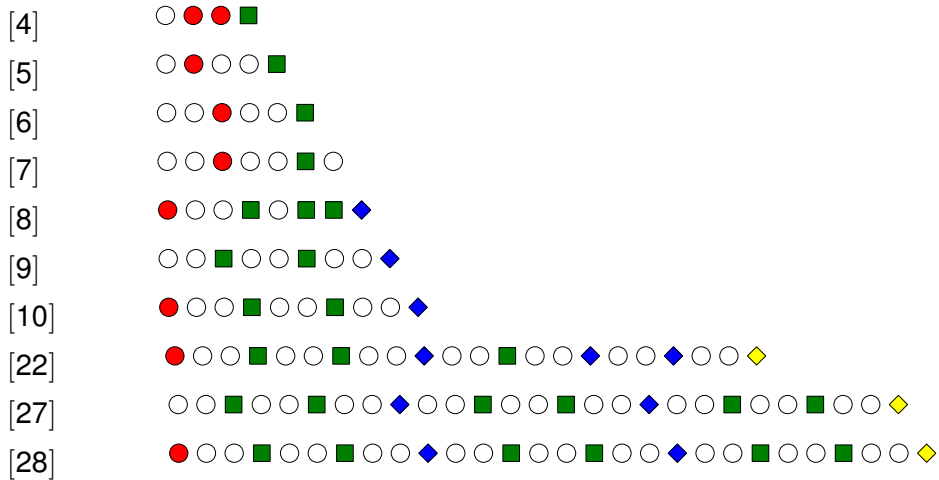
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Monotonicity?

Proposition $\text{aw}([n+1], k) \leq \text{aw}([n], k) + 1.$

Conjecture $\text{aw}([n+1], k) \geq \text{aw}([n], k) - 1.$

Asymptotics of $\text{aw}([n], k)$

Theorem (BEHHKKLMSWY '14)

$$\log_3 n + 2 \leq \text{aw}([n], 3) \leq \log_2 n + 1$$

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For the lower bound, consider the coloring

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$$m_x(3^{c(x)-j}) + m_z(3^{c(z)-j}) = 2m_y(3^{c(y)-j})$$

where m_x, m_y, m_z are relatively prime to 3.

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where m_x, m_y, m_z are relatively prime to 3. Since the colors are distinct, exactly two of the numbers are multiples of three.

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Proposition For $n \geq 2$, there exists $m \leq \lfloor \frac{n}{2} \rfloor$ such that

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There is a minimal interval $[a, b] \subset [n]$ such that all r colors appear. Thus, the color $c(a)$ does not appear within $[a + 1, b]$ and the color $c(b)$ does not appear within $[a, b - 1]$.

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Translate $[a, b]$ to be a coloring of $[t]$ where $t = b - a + 1$.

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Now, $c : [t] \rightarrow [r]$ is an exact coloring where $c(1) \neq c(t)$ and these colors do not appear within $[2, t - 1]$.

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Claim: t is even.

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Claim: t is even. (If not, then $1, \frac{t-1}{2}, t$ is a rainbow 3-AP.)

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Now, $c : [t] \rightarrow [r]$ is an exact coloring where $c(1) \neq c(t)$ and these colors do not appear within $[2, t - 1]$.

Claim: $r - 1$ colors appear on the odd elements of $[t]$.



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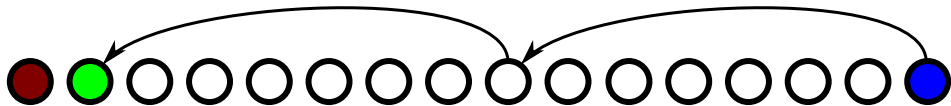
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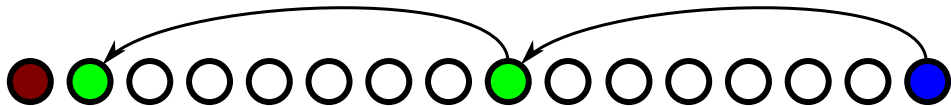
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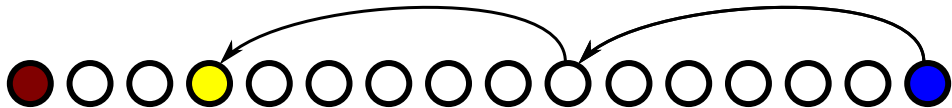
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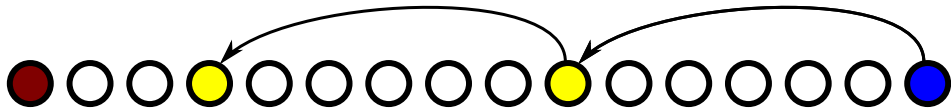
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Note that $r - 1$ colors also appear on the even elements of $[t]$!

Structure of Extremal Colorings



$n = 22$



$n = 28$

The $k \geq 4$ Case

Theorem (BEHHKLSWY '14) For $k \geq 4$,

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Proof.

Let $r + 1 = \text{aw}([n], k)$ and fix an exact r -coloring that avoids rainbow k -APs.

Select one element from each color class. This creates a set S of size r with no k -AP.

$$\text{aw}([n], k) - 1 = |S| \leq \text{sz}([n], k).$$



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This is only non-trivial when $k \geq 6$.

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Let $S \subset [n]$ contain no punctured 4-AP.

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Proposition (BEHHKLSWY '14) Behrend's construction also avoids **punctured 4-APs**: sets given by taking a 4-AP A and removing an element.

Let $S \subset [n]$ contain no punctured 4-AP. If we color S with distinct colors, then $[n] - S$ with a new color, the coloring avoids rainbow 4-APs.

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We can define arithmetic progressions on any additive group, including \mathbb{Z}_n .

Definition $sz(\mathbb{Z}_n, k)$ is the maximum size of a k -AP-free set $S \subset \mathbb{Z}_n$.

Definition $aw(\mathbb{Z}_n, k)$ is the minimum r such that all exact r -colorings of \mathbb{Z}_n contain a rainbow k -AP.

Assume $k \leq n$ and observe:

$$k \leq aw(\mathbb{Z}_n, k) \leq aw([n], k)$$

anti-van der Waerden on \mathbb{Z}_n

	0	1	2	3	4	5	6	7	8	9
0-9				3	3	3	4	3	3	4
10-19	4	3	4	3	4	4	3	4	5	3
20-29	4	4	4	3	4	4	4	5	4	3
30-39	5	4	3	4	5	4	5	3	4	4
40-49	4	4	5	4	4	5	4	3	4	4
50-59	5	5	4	3	6	4	4	4	4	3
60-69	5	3	5	5	3	4	5	3	5	4
70-79	5	3	5	4	4	5	4	4	5	3
80-89	4	6	5	3	5	5	5	4	4	4
90-99	6	4	4	5	4	4	4	4	5	5

Computed values of $\text{aw}(\mathbb{Z}_n, 3)$ for $n = 3, \dots, 99$

The row label gives the range of n and the column heading is the ones digit within this range.

anti-van der Waerden on \mathbb{Z}_n , $k = 3$

Theorem (BEHHKKLMSWY '14)

1. For all positive integers m , $\text{aw}(\mathbb{Z}_{2^m}, 3) = 3$.
2. For an integer $n \geq 2$ having every prime factor less than 100,

$$\text{aw}(\mathbb{Z}_n, 3) = 2 + f_2 + f_3 + 2f_4.$$

Here f_4 denotes the number of odd prime factors of n in the set $Q_4 = \{17, 31, 41, 43, 73, 89, 97\}$. The quantity f_3 is the number of odd prime factors of n in Q_3 , where Q_3 is the set of all odd primes less than 100 and not in Q_4 . Both f_3 and f_4 are counted according to multiplicity. Finally, $f_2 = 0$ if n odd and $f_2 = 1$ if n is even.

anti-van der Waerden on \mathbb{Z}_n , $k \geq 4$

Theorem. For $k \geq 4$,

$$ne^{-O(\sqrt{\log n})} < \text{aw}(\mathbb{Z}_n, k) \leq ne^{-\log \log \log n - \omega(1)}.$$

Proof is essentially the same as the $\text{aw}([n], k)$ case.

Open Problems

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Conjecture For positive integers n and k , $\text{aw}([n], k) \geq \text{aw}([n-1], k) - 1$.

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Conjecture Let m be a nonnegative integer. Then $\text{aw}([3^m], 3) = m + 2$.

Question Is it true that $\text{aw}([3n], 3) = \text{aw}([n], 3) + 1$ for all positive integers n ?

Open Problems

A **singleton extremal coloring** of S is an exact coloring of S that avoids rainbow k -APs and uses exactly $aw(S, k) - 1$ colors

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Conjecture For $k = 3$, there exists a singleton extremal coloring of $[n]$ and of \mathbb{Z}_n .

Conjecture For p an odd prime and $t \geq 3$,

$$\text{aw}(\mathbb{Z}_{pt}, 3) \geq \text{aw}(\mathbb{Z}_t, 3) + \text{aw}(\mathbb{Z}_p, 3) - 2.$$

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Question Does there exist a prime p such that $\text{aw}(\mathbb{Z}_p, 3) \geq 5$?

Rainbow Arithmetic Progressions II: The Collaboration

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May 19th, 2014

University of Colorado Denver

Process

Process is a series of steps or actions that are performed in a specific order to achieve a goal or complete a task. It can be used in various contexts, such as manufacturing, business, and education.

There are several types of processes, including linear processes, cyclical processes, and parallel processes. Each type has its own characteristics and is used in different situations.

Linear processes involve a single sequence of steps that must be completed in order. Cyclical processes involve a series of steps that repeat over and over again. Parallel processes involve multiple steps that can be completed simultaneously.

Understanding the different types of processes and how they are used is essential for managing and improving them. This knowledge can help organizations and individuals to work more efficiently and effectively.

Process management is the practice of identifying, defining, and controlling the processes that are used to create products or services. It is a key component of many business and organizational systems.

Effective process management can lead to improved productivity, reduced costs, and higher quality. It is a continuous process that requires ongoing monitoring and adjustment.

By understanding and managing processes, organizations can ensure that they are always working towards their goals and providing the best possible service to their customers.

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People take turns at board with ideas, and taking suggestions from group.

What Worked

Collaboration on steroids!

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Computation!

A new problem.

What Didn't Work

Big group!

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Ideal: no more than 4 grad students and 2 faculty per group.

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An 11-author paper will look confusing on anyone's C.V.

Ideas for Next Time

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Group Rotation: Have one group meet in seminar room per week. Other groups meet students-only in another room.

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