

## MATH 482, Spring 2013 - Homework 3 Solutions

1. [5pts] Solve (in mixed strategies) for both players the game with the payoff matrix

$$\begin{pmatrix} -2 & 3 & 0 & -6 & 3 \\ 0 & -4 & 9 & 2 & 1 \\ 6 & -2 & 7 & 4 & 5 \\ 7 & -3 & 8 & 3 & 2 \end{pmatrix}.$$

*Solution:* There are multiple ways to arrive at the optimal strategies, but we only need two strategies with value that is feasible in the primal and the dual. However, the solutions **MUST** be presented as fractions, otherwise they are not exact!

$$P_1 = (0, 0, 2/11, 0, 9/11), \quad P_2 = (8/11, 3/11, 0, 0), \quad v_1 = v_2 = 27/11$$

Example Sage Worksheet is online.

One method to find this solution is to reduce the number of strategies using *domination*. If column  $i$  is coordinate-wise smaller than column  $j$ , then Player 1 would never choose strategy  $i$  over strategy  $j$ , since all payoff options are worse for Player 1 in column  $i$ . If row  $a$  is coordinate-wise larger than row  $b$ , then Player 2 would never choose strategy  $a$  over strategy  $b$ , since all payoff options are worse for Player 2 in row  $a$ .

Thus, for this matrix we can delete column 4 (dominated by column 3) and column 2 (dominated by column 4). We can then delete row 3 (dominates row 2) and row 4 (dominates row 1). Finally, we can now delete column 1 (dominated by column 5). This leaves the 2-by-2 matrix  $\begin{bmatrix} 0 & 3 \\ 9 & 1 \end{bmatrix}$  corresponding to the submatrix with rows 1 & 2, and columns 3 & 5. Solving this smaller game is equivalent to the larger game.

2. Add slack variables  $y_1$  &  $y_2$  to be in standard form.  
 Use  $y_1$  &  $y_2$  as initial basis w/  $B=I_2$ ,  $c_B=0$ ,  $x_B=b$ .

CARRY<sup>(1)</sup>:

0	0	0	
5	1	0	$x_1$
8	0	1	$y_2$
			$y_1, y_2$

-7
2
3
$x_1$

CARRY<sup>(2)</sup>:

$35/2$	$7/2$	0	
$5/2$	$1/2$	0	$x_1$
$1/2$	$-3/2$	1	$y_2$
			$y_1, y_2$

$-1/2$
$3/2$
1/2
$x_2$

$$\bar{c}_2 = -11 + \left(\frac{7}{2} \ 0\right) \begin{bmatrix} 3 \\ 5 \end{bmatrix} = -1/2$$

$$X_2 = B^{-1}A_2 = \begin{pmatrix} 1/2 & 0 \\ -3/2 & 1 \end{pmatrix} \begin{bmatrix} 3 \\ 5 \end{bmatrix} = \begin{bmatrix} 3/2 \\ 1/2 \end{bmatrix}$$

CARRY<sup>(3)</sup>:

18	2	1	
1	5	-3	$x_1$
1	-3	2	$x_2$
			$y_1, y_2$

$$\bar{c}_3 = -4 + (2 \ 1) \begin{bmatrix} 1 \\ 2 \end{bmatrix} = 0$$

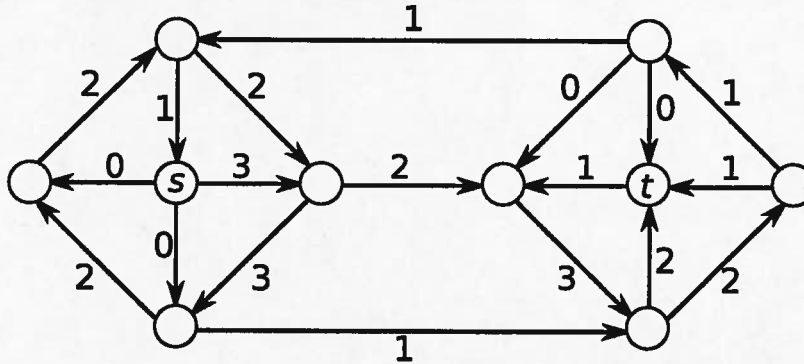
$$\bar{c}_4 = -1 + (2 \ 1) \begin{bmatrix} 2 \\ 1 \end{bmatrix} = 4$$

All relative costs are nonnegative,  
 so optimal!

Solution:  $(1, 1, 0, 0)$   
 $x_1 \ x_2 \ x_3 \ y_1$

with value: -18.

3. [5pts] Let  $G$  be the network with the flow drawn below on the left (the numbers correspond to flow values). Write the flow as a linear combination of flows along cycles and paths from  $s$  to  $t$ .

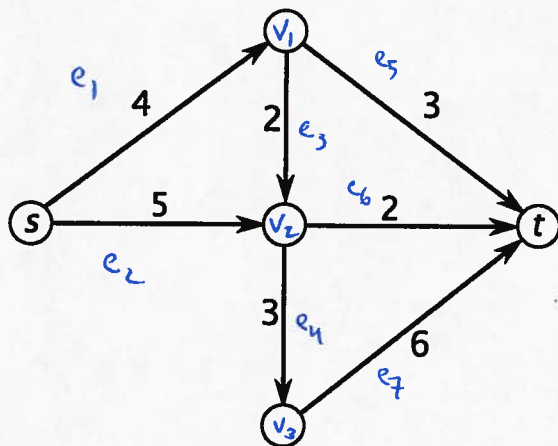


Label the vertices around  $s$  in clockwise order (starting at the top) as  $v_1, v_2, v_3, v_4$ . Label the vertices around  $t$  in clockwise order (starting at the top) as  $v_5, v_6, v_7, v_8$ . Then, the flow can be described as:

$$2 \underbrace{sv_2v_8v_7t}_{\text{path}} + 2 \underbrace{tv_8v_7v_6t}_{\text{cycle}} + 1 \underbrace{sv_2v_3v_7v_6v_5v_1s}_{\text{cycle}} + 2 \underbrace{v_1v_2v_3v_4}_{\text{cycle}}$$

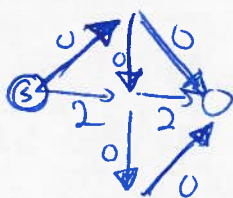
where the flow across an edge is the sum of the flows across all cycles and paths using that edge. (There are multiple ways to solve this problem; all involve at least one cycle and at least one path.)

4. [5pts] Solve the flow problem above and on the right (now below) using revised simplex (the numbers correspond to capacity values).



We start by modeling flow across any st-path.

$$P_1 = s v_2 t, \text{ value} = 2.$$



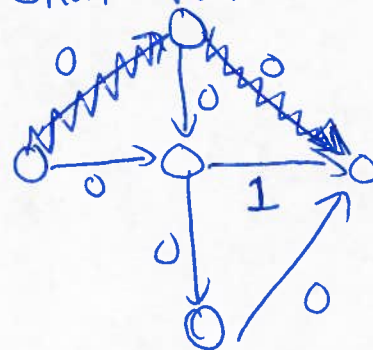
CARRY<sup>(1)</sup>

	0	0	0	0	0	0	0	0	-1
4	1								0
5		1							-1
2			1						0
3				1					0
3					1				0
2						1			-1
6							1		0
	$e_1$	$e_2$	$e_3$	$e_4$	$e_5$	$e_6$	$e_7$		

CARRY<sup>(2)</sup>

	2	1	0	0	0	0	0	1	0	-1
4		1								-1
5			1						-1	0
2				1					0	0
3					1				0	0
3						1			-1	0
2							1		0	0
6								1	0	0
	$f_{P_1}$									$f_{P_2}$

$e_R = -1 + (0000010)$   
Shortest Path!



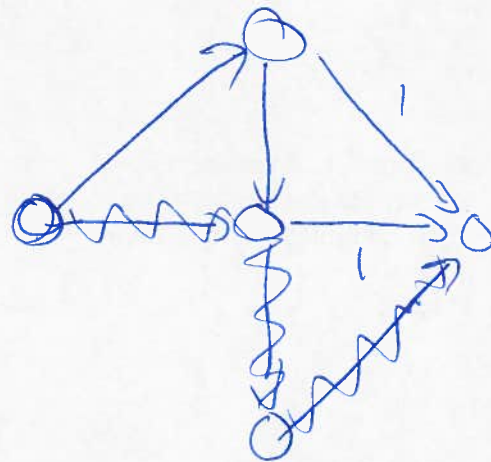
$$P_2 = s v_1 t$$

(4 Continued) CARRY (3)

	5	0	0	0	0	1	1	0
$f_{P_1}$	1	1				-1		
$f_{P_2}$	3		1				-1	
$f_{P_3}$	2			1				
$f_{P_4}$	3				1			
$f_{P_5}$	3					1		
$f_{P_6}$	2						1	
$f_{P_7}$	6							1

-1
0
1
0
1
0
0
0
1

either walk, choose 700



CARRY (4)

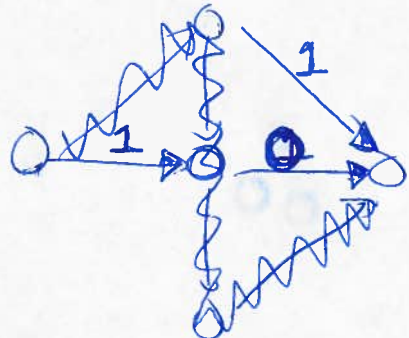
	8	0	1	0	0	1	0	0
$f_{P_1}$	1	1				-1		
$f_{P_2}$	3		1				-1	
$f_{P_3}$	2			1				
$f_{P_4}$	0	-1			1		1	
$f_{P_5}$	3					1		
$f_{P_6}$	2						1	
$f_{P_7}$	3	-1					1	1

-1
1
0
1
0
0
0
1

Observe? Degeneracy!

$$P_3 = 5 v_2 v_3 t$$

(could choose  $P'_3 = 5 v_1 v_2 v_3 t$ )



CARRY (5)

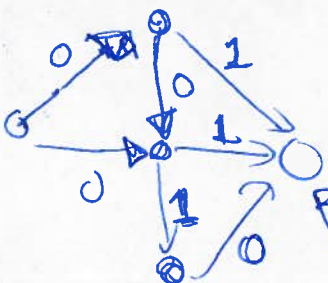
	8	0	0	0	1	1	1	0
$f_{P_1}$	1	1	1		-1	-1	-1	
$f_{P_2}$	3		1				-1	
$f_{P_3}$	2		1	1	-1		-1	
$f_{P_4}$	0	-1			1		1	
$f_{P_5}$	3					1		
$f_{P_6}$	2						1	
$f_{P_7}$	3	0			-1		0	1

$e_1, e_2, e_3, e_4, e_5, e_6, e_7$

$$P_4 = 5 v_1 v_2 v_3 t \quad (\text{could choose } P'_4 = 5 v_1 v_2 t)$$

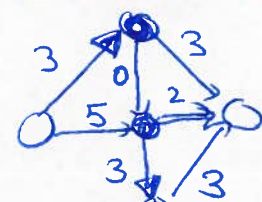
$$\bar{c}_4 = -1 + 0 = -1$$

$$X_{P_4} = B^{-1} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}$$



No path of weight  $< 1$ , so we are optimal!

Edge flows:



Value: 8 (as a flow)