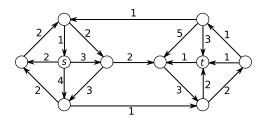
## MATH 482, Spring 2013 - Homework 4 Due Wednesday, 03/27.

Solve 4 of the first 5 problems below, and also solve problem 6. Students registered for 4 credits must solve all problems.

For problems 1–3, consider the following graph.



1.[5pts] (*Primal-Dual Simplex*) Solve the shortest *st*-path problem using the primal-dual simplex algorithm.

**2.**[5pts] (*Dijkstra's Algorithm*) Use Dijkstra's Algorithm to compute the distances from s to all other vertices in the graph above.

**3.**[5pts] (*Max-Flow/Min-Cut*) Solve the max-*st*-flow problem using the Ford-Fulkerson Algorithm. Prove optimality using duality.

**4.**[5pts] (*Floyd-Warshall Algorithm*) Use the Floyd-Warshall Algorithm to compute shortest distances among all pairs of vertices in the graph given by the following adjacency matrix. (The entry  $a_{i,j}$  stores the length of the arc (i, j).)

ΓO	1	2	0	0
0	0	$     \begin{array}{c}       2 \\       0 \\       0 \\       0 \\       0     \end{array} $	4	0
6	1	0	0	3
5	3	0	0	0
0	0	0	3	0

**5.**[5pts] Consider the following linear program.

Solve the linear problem graphically, then also solve the problem graphically when  $x_1$  and  $x_2$  are constrainted to be integers, demonstrating a gap between the real and integer solutions.

**6.**[10pts] (*Matchings and Vertex Covers*) Let G be a graph with edges spanning two sets of vertices, X and Y. A matching is a set M of edges, where  $M = \{x_iy_i : 1 \le i \le k\}$  for some k,  $x_i \in X$ , and  $y_i \in Y$ , with  $x_iy_i \in E(G)$ . A vertex cover is a set  $Q \subset V(G)$  such that all edges have at least one endpoint in Q. Use Max-Flow/Min-Cut to prove that the maximum size of a matching in a bipartite graph G is equal to the minimum size of a vertex cover. (*Hint:* Add vertices s and t to G, direct the edges, and show that the max st-flow and min st-cut problems are equivalent to the max matching and min vertex cover problems.)