

FARKAS LEMMA: (FOR CONES)

LEMMA: (FARKAS)

EITHER

1) $Ax = b, x \geq 0$ HAS SOLUTION
OR

2) $A^T y \geq 0, b^T y < 0$
HAS SOLUTION. BUT NOT BOTH!

REFORMULATION:

$$K = \{Ax : x \in \mathbb{R}_+^n\} \subseteq V = \mathbb{R}^m$$

\rightarrow CLOSED CONVEX CONE

IF $Ax = b$ HAS SOLUTION $\Rightarrow b \in K$

NO SOLUTION $\Rightarrow b \in V \setminus K$

\Rightarrow THM 3.3.2 \Rightarrow

$\Rightarrow \exists y: y^T Ax \geq 0 \forall x \in \mathbb{R}_+^n$ & $y^T b < 0$

$\hat{A}^T \{x \geq 0\} \subseteq \mathbb{R}_+^m$

$$A^T y \in (\mathbb{R}_+^n)^* = \mathbb{R}_+^n$$

(NOT BOTH 1 & 2 SOLVABLE EASY BY $y^T \cdot = 0^T$)

GENERALIZATION:

$$A(x) = \begin{matrix} \uparrow \\ \uparrow \\ \uparrow \end{matrix} \begin{matrix} \text{LINEAR OPERATOR} \\ W \\ \text{SOME CONE} \end{matrix} \quad x \in K \subseteq V$$

LINEAR OPERATOR $V \rightarrow W$

PROBLEMS:

- A^T ?
- $\{A(x), x \in K\}$ NOT CLOSED
 \uparrow CONE CONE

DEF:

$A: V \rightarrow W$ LINEAR OPERATOR.

LINEAR OPERATOR $A^T: W \rightarrow V$ IS

ADJOINT OF A IF $\forall x \in V, \forall y \in W$

$$\langle y | A(x) \rangle = \langle A^T(y) | x \rangle$$

NOTE:

- IF A MATRIX, ADJOINT IS A^T
- IF A^T EXISTS, UNIQUE
- IF V, W FINITE DIMENSION, A^T EXISTS

FARKAS LEMMA 1ST TRY

LET $K \subseteq V$ CONE, $b \in W$

EITHER

1) $A(x) = b, x \in K$ HAS SOLUTION

OR

2) $A^T(y) \in K^\circ, \langle b, y \rangle < 0, y \in W$
HAS SOLUTION BUT NOT BOTH.

NOT BOTH SOLVABLE TOGETHER: IF $A(x) = b$ THEN

$$0 > \langle b, y \rangle = \langle y, Ax \rangle = \langle A^T(y), x \rangle \geq 0$$

AT LEAST ONE IF $C = \{A(x), x \in K\}$ CLOSED
HAS SOLUTION \rightarrow (THM 3.2 MAKES IT)

EX) $K = \langle 0 \rangle \subseteq \mathbb{R}^3 = \{x \geq 0, y \geq 0, x + y \geq 1\}$

LET $A = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}$

$$C = \{(x, t) \in \mathbb{R}^2 : (x, y, t) \in K\} = 0 \cup \{x > 0\} \times \mathbb{R}$$

NOT CLOSED

FOR $b = (0, 1)$, NONE OF 1, 2, SOLVABLE.

1) CLEAR) $y = (1, -1)$ $\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \rightarrow 0 \Rightarrow$

2) $A^T y \in \langle 0 \rangle$ SEE $\langle 0 \rangle$ NO SOLUTION

FARKAS LEMMA 2nd TRY

LEMMA:

$K \subseteq V$ CLOSED CONVEX CONE,

$$C = \{A(x) : x \in K\}$$

CLOSURE OF $C = \bar{C}$ IS CLOSED CONVEX CONE

DEFINITION:

$A(x) = b, x \in K$ IS SUBFEASIBLE

IF $\exists (x_k)_{k \in \mathbb{N}} \subseteq K$ ST. $\lim_{k \rightarrow \infty} A(x_k) = b$
(THEN $b \in \bar{C}$)

LEMMA (FARKAS FOR CONES)

EITHER $A(x) = b, x \in K$ IS SUBFEASIBLE
OR $A^T(\gamma) \in K^\circ, \langle b, \gamma \rangle < 0$ HAS SOLUTION

PROOF:

$\rightarrow A(x) = b, x \in K$ IS SUBFEASIBLE: $(x_k)_{k \in \mathbb{N}}$ IS SEQ

$$\text{FOR } \gamma: \langle \gamma, b \rangle = \langle \gamma, \lim_{k \rightarrow \infty} A(x_k) \rangle = \\ = \lim_{k \rightarrow \infty} \langle \gamma, A(x_k) \rangle = \lim_{k \rightarrow \infty} \langle A^T(\gamma), x_k \rangle \geq 0$$

\Rightarrow OR HAS NO SOLUTION. IF $A^T(\gamma) \in K^\circ$

• $A(x) = b$ NOT SUBFEASIBLE - FROM SEPARATION THM 0

CONE PROGRAMS

DEFINITION:

$K \subseteq V, L \subseteq W$ CLOSED CONVEX CONES,
 $b \in W, c \in V, A: V \rightarrow W$ LINEAR OP.

CONE PROGRAMS:

$$(P) \begin{cases} \max & \langle c, x \rangle \\ \text{s.t.} & b - A(x) \in L \\ & x \in K \end{cases}$$

OR IF $L = \{0\}$, WE GET EQUATIONS

• FEASIBLE: $\exists x \in K$ S.T. $b - A(x) \in L$

• VALUE OF (P)

$$\sup \{ \langle c, x \rangle, x \text{ FEASIBLE} \} = \infty \quad \text{MAYBE}$$

• x^* OPTIMAL IF $\langle c, x^* \rangle \geq \langle c, x \rangle \forall x \text{ FEASIBLE}$

⊗ UNOBTAINABLE FINITE VALUE

$$\begin{array}{ll} \min x & z=1 \\ \text{s.t. } z=1 & \text{subst.} \\ (x, z) \in C & \rightarrow \\ & x \geq 0 \end{array}$$

VALUE = 0 BUT
NOT OBTAINABLE. NO OPT SOL.

AVOID TROUBLES WHEN (P) NOT FEASIBLE BUT
TINY PERTURBATION MAKES (P) FEASIBLE
(NOT THE CASE FOR (LP))

DEF (SUBFEASIBILITY)

(P) IS SUBFEASIBLE IF $\exists (x_k)_{k \in \mathbb{N}} \subseteq K$
AND $\exists (x'_k)_{k \in \mathbb{N}} \subseteq L$ SUCH THAT
 $\lim_{k \rightarrow \infty} (Ax_k + x'_k) = b$

x_k, x'_k ARE SUBFEASIBLE SEQUENCES

VALUE OF SUBFEASIBLE SEQUENCE:

$$\langle c, (x_k) \rangle := \lim_{k \rightarrow \infty} \langle c, x_k \rangle$$

SUBVALUE OF (P) = $\max_{\text{FEASIBLE SEQ.}} \{ \langle c, (x_k) \rangle \}$

! VALUE OF (P) \leq SUBVALUE OF (P), $\left\{ \begin{array}{l} \text{EX } \Delta > 0 \\ \dots \end{array} \right.$
(EVEN IF BOTH $\uparrow \uparrow$ FINITE, STILL Δ CAN BE)

DEF: $x \in K$ IS INTERIOR POINT IF $\exists A(x) \in \text{INT}(L)$

THEOREM: (P) HAS INTERIOR POINT \Rightarrow

\Rightarrow VALUE = SUBVALUE

DUAL PROGRAM

$$(D) \begin{cases} \text{MIN } \langle l, y \rangle \\ \text{s.t. } A^T(y) - c \in K^* \\ y \in L^* \end{cases}$$

∴ DUAL OF (D) IS (P)

SERIES OF THEOREMS:

THEOREM: WEAK DUALITY

IF (D) FEASIBLE & (P) SUBFEASIBLE \Rightarrow
SUBVALUE(P) \leq VALUE(D) \leftarrow EASY

THEOREM: REGULAR DUALITY

(D) IS FEASIBLE & HAS FINITE VALUE $\beta \Rightarrow$
 \Rightarrow (P) IS SUBFEASIBLE & HAS FINITE
SUBVALUE γ , MORE OVER, $\gamma = \beta$

\hookrightarrow DONE FOR LP, FROM FARKAS LEMMA

BUT STILL VALUE(P) SUBVALUE(P)



SUBVALUE(D) \uparrow VALUE(D)

GAP = - ϵ

THEOREM STRONG DUALITY

IF (P) FEASIBLE, FINITE VALUE γ , AND HAS INTERIOR POINT THEN (D) IS

FEASIBLE AND HAS FINITE VALUE $\beta = \gamma$.

(INT POINT \Rightarrow VALUE = DUAL VALUE)

FOR SDP:

NOTE: IF $L = \{\emptyset\}$ DOES NOT APPLY \because NO INT.

THEOREM:

IF $\left\{ \begin{array}{l} \text{MAX } \langle c, x \rangle \\ \text{ST } A(x) = b \\ x \in K \end{array} \right.$

(P) $\left\{ \begin{array}{l} \text{ST } A(x) = b \\ x \in K \end{array} \right.$

← SATISFY

IS FEASIBLE, HAS FINITE VALUE γ , $x \in \text{INT}(K)$

THEN $\text{MIN } \langle b, y \rangle$

ST. $A^T(y) - c \in K^*$

IS FEASIBLE AND HAS FINITE VALUE $\beta = \gamma$

LEMMA: $(\text{PSD}_n)^* = \text{PSD}_n$

EXERCISE

\Rightarrow