

## Linear Programming Algorithms - Interior point methods

Source: Chapter 11 of Convex Optimization, Stephen Boyd and Lieven Vandenberghe

Let

$$(P) \begin{cases} \text{minimize} & f(\mathbf{x}) \\ \text{subject to} & g_i(\mathbf{x}) \leq 0 \text{ for } i = 1, \dots, m \end{cases}$$

where  $f, g$  are convex, twice continuously differentiable and optimal solution  $\mathbf{x}^*$  exists. Moreover, let  $(P)$  be *superconsistent*,  $\exists \mathbf{x}, \forall i, g_i(\mathbf{x}) < 0$ .

(the setup covers linear, quadratic, geometric, semidefinite, ... programming).

**Idea:** Change the  $(P)$  to a problem without constraints.

Let

$$(P') = \text{minimize } f(\mathbf{x}) + \sum_{i=1}^m I(g_i(\mathbf{x})),$$

where  $I$  is an indicator function

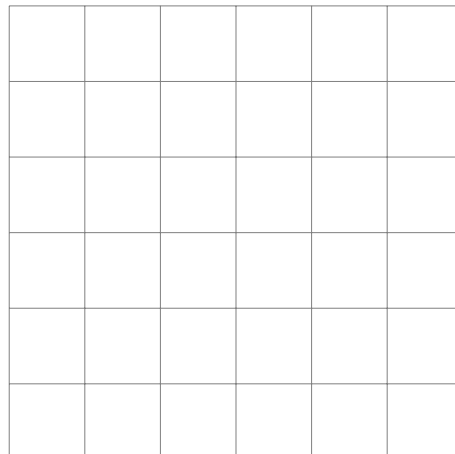
$$I(u) = \begin{cases} 0 & \text{if } u \leq 0 \\ +\infty & \text{if } u > 0. \end{cases}$$

**1:** What is the optimal solution to  $(P')$ ?

**2:** Can you solve  $(P')$  by methods from calculus?

Use approximation of  $I(u) \approx -c \log(-u)$ , where  $c > 0$ .

**3:** Sketch  $I(u)$  and its approximations. Is the approximation better when  $c$  is large or small?



For  $t > 0$ , we consider a smooth unconstrained approximation of  $(P')$

$$\text{minimize } f(\mathbf{x}) - \frac{1}{t} \sum_{i=1}^m \log(-g_i(\mathbf{x})).$$

Define *logarithmic barrier function*

$$\Phi(\mathbf{x}) = - \sum_{i=1}^m \log(-g_i(\mathbf{x})),$$

for all  $\mathbf{x}$  where  $g_i(\mathbf{x}) < 0$  (interior of feasible solutions)

**Analytic center** of the set  $S = \{\mathbf{x} : g_i(\mathbf{x}) \leq 0\} \subseteq \mathbb{R}^n$  is  $\mathbf{x}^*$  minimizing  $\Phi(\mathbf{x})$  over all  $\mathbf{x} \in S$ .

**4:** Find the analytic center of a square in  $\mathbf{R}^2$  defined by equations

$$x_1 \geq 0, x_2 \geq 0, x_1 \leq 1, x_2 \leq 1$$

**5:** Find the analytic center of a square in  $\mathbf{R}^2$  defined by equations

$$x_1 \geq 0, x_2 \geq 0, (1 - x_1)^3 \geq 0, (1 - x_2)^3 \geq 0.$$

Notice it is possible to define center even if functions are not convex everywhere and the center depends on the functions.

For  $t > 0$  define  $\mathbf{x}^*(t)$  as the optimal solution of

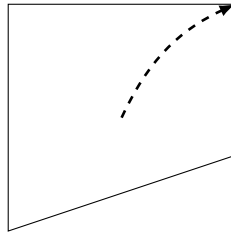
$$(P_t) = \text{minimize } tf(\mathbf{x}) + \Phi(\mathbf{x}).$$

(assume that the optimal solution is unique)

**Central path** is  $\{\mathbf{x}^*(t) : t \geq 0\}$ .

Interior point method idea: Start in the analytical center and follow the central path.

In iterations increase  $t$  and recompute the new optimum by Newton's method.



There exists a notion of dual program ( $D$ ) for ( $P$ ), (based on Karush-Kuhn-Tucker theorem). It gives solutions to the dual  $\mathbf{y}^*(t)$  such that

$$f(\mathbf{x}^*(t)) - h(\mathbf{y}^*(t)) \leq \frac{m}{t}.$$

Hence the central path converges to  $\mathbf{x}^*$  for ( $P$ ).

**6:** Let  $(LP) = \text{minimize } \mathbf{c}^T \mathbf{x}$  s.t.  $A\mathbf{x} \leq \mathbf{b}$ , where  $A \in \mathbb{R}^{m \times n}$ . Write  $\Phi(\mathbf{x})$ . Use tools from calculate conditions for the corresponding central path. What can you say about the Hessian?

In the end we obtained

$$t\mathbf{c} = -\nabla\Phi(\mathbf{x}^*(t)).$$

Since  $\nabla\Phi(\mathbf{x})$  is perpendicular to the level curve  $\{\mathbf{x} : \Phi(\mathbf{x}) = \Phi(\mathbf{x}^*(t))\}$ , plane  $\mathbf{c}^T \mathbf{x} = \mathbf{c}^T \mathbf{x}^*(t)$  is a tangent of the level curve for  $\Phi$ .

