

## Turán's Theorem

**Theorem 1** (weak Turán's Theorem, 1941). *The maximum number of edges in a graph on  $n$  vertices with no  $(k+1)$ -clique subgraph is at most*

$$\left(1 - \frac{1}{k}\right) \frac{n^2}{2}.$$

Let  $T_k(n)$  be a complete  $k$ -partite graph on  $n$  vertices with parts of as equal sizes as possible, i.e., sizes are  $\lfloor \frac{n}{k} \rfloor$  or  $\lceil \frac{n}{k} \rceil$ . Another way of defining it would be a balanced blow-up of  $K_k$ . Such graph is called the *Turán graph*.

**1:** Show that  $T_k(n)$  gives asymptotically tight lower bound for Turán's theorem.

**Solution:** It is easy to see that  $T_k(n)$  cannot contain a  $(k+1)$ -clique as any set of  $k+1$  vertices in  $T_k(n)$  will have two vertices in the same class and therefore not connected by an edge. Furthermore,  $e(T_k(n)) \geq \binom{k}{2} (\lfloor \frac{n}{k} \rfloor)^2 > \binom{k}{2} (\frac{n}{k} - 1)^2 = (1 - \frac{1}{k}) \frac{n^2}{2} - O(n)$ .

**2:** Prove Turán's theorem by induction. Idea: Find a clique  $A$  of size  $k$  and remove it for induction.

**Solution:** [First proof.] (Turán, 1941) Induction on  $n$ . The theorem is trivially true for  $n \leq k$ , so let  $n > k$  and assume the theorem holds for smaller graphs. Let  $G$  be a graph on  $n$  vertices with no  $(k+1)$ -clique and the maximum number of edges. Therefore,  $G$  must contain a  $k$ -clique as otherwise we could add edges to  $G$  contradicting maximality. Let  $A$  be a clique of size  $k$  and let  $B$  be the remaining  $n - k$  vertices. The graph  $B$  has no  $(k+1)$ -clique so by induction  $e(B) \leq (1 - \frac{1}{k}) \frac{(n-k)^2}{2}$ . Furthermore, each vertex in  $B$  can have at most  $k - 1$  neighbors in  $A$ , so we have

$$e(G) \leq \binom{k}{2} + \left(1 - \frac{1}{k}\right) \frac{(n-k)^2}{2} + (n-k)(k-1) = \left(1 - \frac{1}{k}\right) \frac{n^2}{2}.$$

Actually, Turán proved more.

**Theorem 2** (Turán's Theorem, 1941). *The maximum number of edges in an  $n$ -vertex graph with no  $(k+1)$ -clique is exactly  $e(T_k(n))$ . Furthermore,  $T_k(n)$  is the unique graph attaining this maximum.*

We show several more proofs that prove one version or the other. Note that  $T_k(n)$  is a complete multipartite graph and among complete multipartite graphs with no  $(k+1)$ -clique is it the largest. This leads to an ingenious approach: if we can show that a graph  $G$  with no  $(k+1)$ -clique and the maximum number of edges is complete multipartite then we are done.

**3:** Prove  $T_k(n)$  is the unique extremal graph by using induction on  $k$ . Idea: Take a vertex  $x$  of maximum degree, use induction on neighbors of  $x$ , and maximality arguments on non-neighbors of  $x$ .

**Solution:** [Second proof.] (Erdős, 1970) Induction on  $k$ . The theorem is trivially true for  $k = 1$ , so let  $k > 1$  and assume the theorem holds for  $k - 1$ . We will prove that if  $G$  has no  $(k+1)$ -clique and the maximum number of edges, then  $G = T_k(n)$ . Let  $x$  be a vertex of maximum degree, let  $S$  be the neighbors of  $x$  and let  $T = G - S$  be the remaining vertices. The graph  $S$  has no  $k$ -clique (as otherwise we could build a  $(k+1)$ -clique with  $x$ ). Let us construct a new graph  $H$  on the vertex set of  $G$  as follows. The graph  $H$  is the same as  $G$  on  $S$ , it contains all edges between  $S$  and  $T$ ,

and it has no edges in  $T$ . Observe that all degrees in  $H$  are at least as large as in  $G$ . Thus  $H$  has at least as many edges as  $G$ . However, if the graph  $G$  has an edge in  $T$ , then  $H$  would have more edges. Therefore, as  $G$  was maximal,  $T$  contains no edges in  $G$  and thus  $G = H$ . Because  $S$  does not contain a  $k$ -clique,  $H$  does not contain a  $(k + 1)$ -clique and therefore  $H$  is a graph with the maximum number of edges with this property. Furthermore,  $S$  must be an edge-maximal graph with no  $k$ -clique. Thus, by induction on  $k$ , we have  $S = T_{k-1}(|S|)$ , so  $H$  is a complete multipartite graph. The largest (most edges)  $n$ -vertex complete multipartite graph with no  $(k + 1)$ -clique is  $T_k(n)$ .

Another proof is based on the characterization of complete multipartite-graphs by a forbidden subgraph on 3 vertices.

**4:** Let  $G$  be a graph with no three vertices that induce exactly one edge. That is,  $G$  does not contain  $\overline{P_3}$ , so called co-cherry. Show that  $G$  is a complete multipartite graph.

**Solution:** Suppose  $G$  does not contain  $\overline{P_3}$ . If  $G$  is a complete graph, we are done. Let  $x, y$  be two non-adjacent vertices. Since  $G$  does not contain  $\overline{P_3}$ , all other vertices are partitioned to sets  $A$  and  $B$ , where  $A$  are neighbors of both  $x$  and  $y$ , while  $B$  are non-neighbors of both  $x$  and  $y$ . Let  $a \in A$  and  $b \in B$ . By considering triple  $a, b, x$ , we conclude  $ab$  is an edge. Let  $b_1, b_2 \in B$ . By considering triple  $b_1, b_2, x$ , we conclude  $b_1, b_2$  is not an edge. Hence  $B \cup \{x, y\}$  is an independent set and there is a complete bipartite graph between  $B \cup \{x, y\}$  and  $A$ . Hence  $B \cup \{x, y\}$  is one part of the multi-partite graph and other parts can be obtained by repeating the same argument on  $A$ .

**5:** Prove  $T_k(n)$  is the unique extremal graph by showing the extremal graph must be a complete multipartite graph. Idea: Take a co-cherry and consider degrees of its vertices. Duplicate/erase vertices to get more edges.

**Solution:** [Third proof.] (Zykov, 1949)<sup>1</sup> Let  $G$  be an  $n$ -vertex graph with no  $(k + 1)$ -clique and the maximum number of edges. We will show that  $G$  is a complete multipartite graph. If  $G$  is not multipartite, then there is a pair of non-adjacent vertices  $x$  and  $y$  and a vertex  $z$  such that  $xz$  is an edge, but  $yz$  is not an edge. If  $d(x) > d(y)$ , then remove all edges incident to  $y$  and connect  $y$  to all neighbors of  $x$  (but not  $x$  itself). The resulting graph has no  $(k + 1)$ -clique but has more edges than  $G$  which contradicting maximality. Thus  $d(y) \geq d(x)$  and  $d(y) \geq d(z)$  (by the same argument for  $z$ ). Now remove all edges incident to  $x$  and  $z$  and connect both vertices to the neighbors of  $y$ . The resulting graph has more edges than  $G$  and has no  $(k + 1)$ -clique again contradicting maximality. Therefore  $G$  must be a complete multipartite graph. The largest complete multipartite graph with no  $(k + 1)$ -clique is  $T_k(n)$ .

<sup>1</sup>Due to WWII, Zykov's work was done without the knowledge of Turán's.

[Fourth proof on Turán's Theorem](Katona-Nemetz-Simonovits, 1965)<sup>2</sup>

**6:** Denote the number of edges of the Turán graph on  $n$  vertices with  $k$  classes by

$$t_k(n) = |E(T_k(n))|.$$

Show a technical lemma that

$$(n+1)t_k(n) < (n-1)t_k(n+1) + (n-1)$$

Idea: Write  $n = qk + r$  where  $r < k$  and precisely calculate  $2t_k(n)$  and  $t_k(n) = t_k(n+1) - \text{something}$ . Then prove the inequality.

**Solution:** Put  $n = qk + r$  where  $r < k$ . Because  $n > k$  we have either  $r > 0$  or  $q > 1$ . Then  $T_k(n)$  has  $k - r$  classes of size  $q$  and  $r$  classes of size  $q + 1$ . The maximum degree in  $T_k(n)$  is  $n - q$  and there are  $r(q + 1)$  many vertices of degree 1 less, so summing degrees gives

$$2t_k(n) = n(n - q) - r(q + 1).$$

Adding a vertex to a class of size  $q$  gives  $T_k(n + 1)$ , thus

$$t_k(n) = t_k(n + 1) - (n - q).$$

Now we use the above two equalities to get

$$(n+1)t_k(n) = (n-1)t_k(n) + 2t_k(n) = (n-1)t_k(n+1) + (n-q) - r(q+1).$$

Observe that  $(n - q) - r(q + 1) < n - 1$  as long as either  $r > 0$  or  $q > 1$  which completes the proof.

**7:** Now for the main proof, let  $g_k(n)$  be the maximum number of edges possible in an  $n$ -vertex graph with no  $K_{k+1}$  subgraph. Obviously  $t_k(n) \leq g_k(n)$  we would like to show  $t_k(n) \geq g_k(n)$ .

Idea: Proceed by induction on  $n$ . Double count pairs  $(v, e)$  where  $v$  is a vertex and  $e$  is an edge NOT incident to  $v$ . Use the previous inequality.

**Solution:** We proceed by induction on  $n$ . The assertion is true for  $n \leq k + 1$  so let us assume the inequality holds for  $n \geq k + 1$  and now show it holds for  $n + 1$ .

Let  $G$  be an  $(n + 1)$ -vertex  $K_{k+1}$ -free graph with the maximum number of edges. We will double count the pair  $(v, e)$  where  $v$  is a vertex and  $e$  is an edge NOT incident to  $v$ . We can fix  $v$  first in  $n + 1$  ways and then fix  $e$  in at most  $g_k(n)$  ways. Alternatively, we can fix  $e$  in  $g_k(n + 1)$  ways and then fix  $v$  in exactly  $n - 1$  ways. Combining these estimates and applying induction to  $g_k(n)$  gives

$$g_k(n+1)(n-1) \leq (n+1)g_k(n) \leq (n+1)t_k(n). \quad (1)$$

To complete the proof it is enough to recall that  $(n+1)t_k(n) < (n-1)t_k(n+1) + (n-1)$ .

<sup>2</sup>This proof is from the first paper of Nemetz and Simonovits and the second of Katona.