

FLAG ALGEBRAS AND ITS APPLICATION

Bernard Lidický



CombinaTexas 2023

Apr 23, 2023



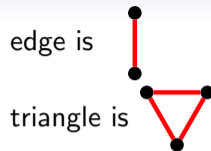
OUTLINE

- Flag Algebras Introduction
- “Proof” of Mantel’s theorem
- Erdős Pentagon Problem and inducibility
- Crossing numbers
- ℓ_2 -norm in Turán type problems
- ε -similar triangles
- Small Ramsey numbers

We will not make it all the way....

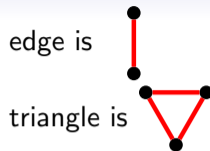
INSPIRATIONAL PROBLEM

- Let n be fixed number of vertices in a graph G .
- Assume G has m edges.
- What is the number of triangles in G ?



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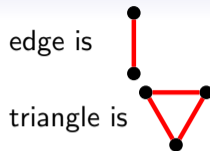
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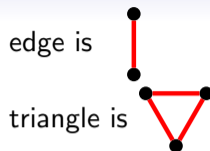
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Consider $n \rightarrow \infty$.

Edges = $p \binom{n}{2}$

Triangles = $t \binom{n}{3}$

Now $p, t \in [0, 1]$.



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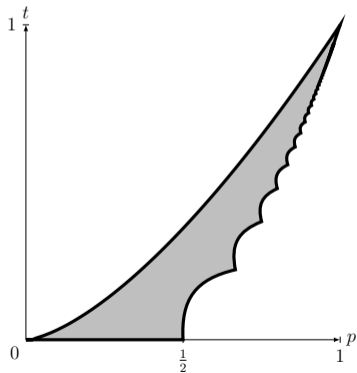
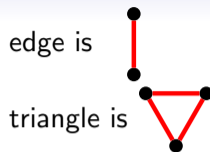
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Upper bound $p^{3/2}$ Kruskal-Katona 1964

Asymptotic lower bound by Razborov 2008



FLAG ALGEBRAS

Seminal paper:

Razborov, Flag Algebras, *Journal of Symbolic Logic* **72** (2007), 1239–1282.

David P. Robbins Prize by AMS for Razborov in 2013 over 300 citations (on google)



EXAMPLE

If density of edges is p , what is the minimum density of triangles?

- Designed to attack extremal problems.
- Works well if constraints as well as desired value can be computed by checking small subgraphs (or average over small subgraphs).
- The results are for the limit as graphs get very large.

EXAMPLE EXTREMAL PROBLEM

THEOREM (MANTEL 1907)

Every n -vertex triangle-free graph contains at most $\frac{1}{4}n^2$ edges.



PROBLEM

Maximize a graph parameter (# of edges) over a class of graphs (triangle-free).

- local condition and global parameter (computable locally)
- threshold
- bound and extremal example

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We will use colors for **edges** and **non-edges**.

FLAG ALGEBRAS DEFINITIONS

Let G be a 2-edge-colored complete graph on n vertices.



The probability that three random vertices in G span a red triangle, i.e. $\# \text{red triangle} / \binom{n}{3}$.

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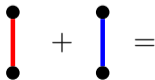
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$$\text{red edge} + \text{blue edge} = 1$$

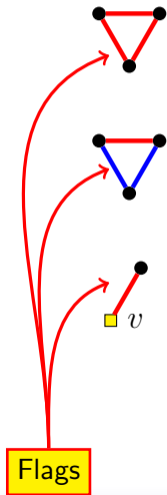
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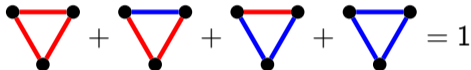
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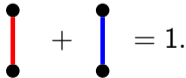
FLAG ALGEBRAS IDENTITIES

Let G be a 2-edge-colored complete graph on n vertices.



The diagram shows four triangles, each with three vertices and three edges. The first triangle has all three edges colored red. The second triangle has the top edge colored blue and the two bottom edges colored red. The third triangle has the top edge colored red and the two bottom edges colored blue. The fourth triangle has all three edges colored blue. These four triangles are separated by plus signs, followed by an equals sign and the number 1.

Same kind as



The diagram shows two vertical edges, each with two vertices. The first edge is colored red and the second edge is colored blue. They are separated by a plus sign, followed by an equals sign and the number 1.

FLAG ALGEBRAS IDENTITIES

Let G be a 2-edge-colored complete graph on n vertices.

$$\begin{array}{c} \bullet \\ | \\ \bullet \end{array} = \frac{3}{3} \begin{array}{c} \bullet & \bullet \\ / \quad \backslash \\ \bullet \end{array} + \frac{2}{3} \begin{array}{c} \bullet & \bullet \\ / \quad \backslash \\ \bullet \end{array} + \frac{1}{3} \begin{array}{c} \bullet & \bullet \\ / \quad \backslash \\ \bullet \end{array} + \frac{0}{3} \begin{array}{c} \bullet & \bullet \\ / \quad \backslash \\ \bullet \end{array}$$

Expanded version:

$$P \left(\begin{array}{c} \bullet \\ | \\ \bullet \end{array} \right) = P \left(\begin{array}{c} \bullet \\ | \\ \bullet \end{array} \mid \begin{array}{c} \bullet & \bullet \\ / \quad \backslash \\ \bullet \end{array} \right) \cdot P \left(\begin{array}{c} \bullet & \bullet \\ / \quad \backslash \\ \bullet \end{array} \right) + P \left(\begin{array}{c} \bullet \\ | \\ \bullet \end{array} \mid \begin{array}{c} \bullet & \bullet \\ / \quad \backslash \\ \bullet \end{array} \right) \cdot P \left(\begin{array}{c} \bullet & \bullet \\ / \quad \backslash \\ \bullet \end{array} \right) + \dots$$

Law of total probability

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Law of total probability

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Let G be a 2-edge-colored complete graph on n vertices.

$$\begin{array}{c} \bullet \\ \diagdown \\ \square v \\ \diagup \\ \bullet \end{array} \times \begin{array}{c} \bullet \\ \diagdown \\ \square v \\ \diagup \\ \bullet \end{array} = \begin{array}{c} \bullet \quad ? \quad \bullet \\ \diagdown \quad \diagup \\ \square v \end{array} + o(1) = \begin{array}{c} \bullet \quad \bullet \\ \diagdown \quad \diagup \\ \square v \end{array} + \begin{array}{c} \bullet \quad \bullet \\ \diagdown \quad \diagup \\ \square v \end{array} + o(1)$$

$o(1)$ as $n \rightarrow \infty$ (will be omitted on next slides)

FLAG ALGEBRAS IDENTITIES

Let G be a 2-edge-colored complete graph on n vertices.

$$\begin{array}{c} \bullet \\ \text{red} \\ \square v \end{array} \times \begin{array}{c} \bullet \\ \text{red} \\ \square v \end{array} = \begin{array}{c} \bullet \text{ ? } \bullet \\ \text{red} \quad \text{red} \\ \square v \end{array} + o(1) = \begin{array}{c} \bullet \text{ red } \bullet \\ \text{red} \quad \text{red} \\ \square v \end{array} + \begin{array}{c} \bullet \text{ blue } \bullet \\ \text{red} \quad \text{red} \\ \square v \end{array} + o(1)$$



: The probability of choosing two **different** vertices ...

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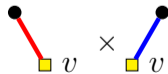
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: The probability of choosing two **different** vertices ...



: The probability that choosing two vertices u_1, u_2 other than v gives red vu_1 and blue vu_2 .

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

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Let G be a 2-edge-colored complete graph on n vertices.

$$\frac{1}{3} \text{ (triangle with blue top edge, red bottom edges) } = \frac{1}{n} \sum_{v \in V(G)} \text{ (triangle with blue top edge, red bottom edges, and a yellow square at the bottom vertex labeled } v \text{)}$$

FLAG ALGEBRAS IDENTITIES

Let G be a 2-edge-colored complete graph on n vertices.

$$\frac{1}{3} \text{triangle} = \frac{1}{n} \sum_{v \in V(G)} \text{triangle}_v$$


$$\text{triangle} \binom{n}{3} = \sum_{v \in V(G)} \text{triangle}_v \binom{n-1}{2}$$


FLAG ALGEBRAS IDENTITIES

Let G be a 2-edge-colored complete graph on n vertices.

$$\frac{1}{3} \text{ (triangle with 2 blue edges) } = \frac{1}{n} \sum_{v \in V(G)} \text{ (triangle with 2 blue edges and vertex } v \text{ highlighted)}$$

$$\text{ (triangle with 3 red edges) } = \frac{1}{n} \sum_{v \in V(G)} \text{ (triangle with 3 red edges and vertex } v \text{ highlighted)}$$

$$\text{ (triangle with 2 blue edges) } \binom{n}{3} = \sum_{v \in V(G)} \text{ (triangle with 2 blue edges and vertex } v \text{ highlighted) } \binom{n-1}{2}$$

FLAG ALGEBRAS IDENTITIES

Let G be a 2-edge-colored complete graph on n vertices.

$$\frac{1}{3} \text{ (triangle with 1 blue edge) } = \frac{1}{n} \sum_{v \in V(G)} \text{ (triangle with 1 blue edge and vertex } v \text{ highlighted)}$$
$$\text{ (triangle with 1 blue edge) } \binom{n}{3} = \sum_{v \in V(G)} \text{ (triangle with 1 blue edge and vertex } v \text{ highlighted) } \binom{n-1}{2}$$

$$\text{ (triangle with 0 blue edges) } = \frac{1}{n} \sum_{v \in V(G)} \text{ (triangle with 0 blue edges and vertex } v \text{ highlighted)}$$
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IDENTITIES SUMMARY

$$1 = \begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagdown \quad \diagup \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagup \quad \diagdown \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagdown \quad \diagup \\ \bullet \end{array}$$

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$$\frac{1}{3} \begin{array}{c} \bullet \\ \diagdown \quad \diagup \\ \bullet \end{array} = \frac{1}{n} \sum_{v \in V(G)} \begin{array}{c} \bullet \\ \diagdown \quad \diagup \\ \square v \end{array}$$

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EXAMPLE - MANTEL'S THEOREM

THEOREM (MANTEL 1907)

A triangle-free n -vertex graph contains at most $\frac{1}{4}n^2 \approx \frac{1}{2} \binom{n}{2}$ edges.



Assume **edges are red** and **non-edges are blue**.

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$$0 \leq \left(1 - 2 \begin{array}{c} \bullet \\ | \\ \square v \end{array} \right)^2$$





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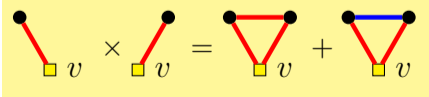
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Assume edges are red and non-edges are blue.

Assume  = 0. (We want to conclude  $\leq \frac{1}{2}$.)

$$0 \leq \left(1 - 2 \text{  } v\right)^2 = \left(1 - 4 \text{  } v + 4 \text{  } v + 4 \text{  } v\right)$$





The diagram shows the multiplication of two edge diagrams (each consisting of a vertex v in a yellow square connected to two vertices in black circles by red edges) resulting in the sum of two triangle diagrams. The first triangle has all three edges in red, and the second triangle has two edges in red and one edge in blue.





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Assume edges are red and non-edges are blue.

Assume  = 0. (We want to conclude  $\leq \frac{1}{2}$.)

$$\begin{aligned}
 0 &\leq \frac{1}{n} \sum_v \left(1 - 2 \cdot \text{edge}(v) \right)^2 = \frac{1}{n} \sum_v \left(1 - 4 \cdot \text{edge}(v) + 4 \cdot \text{triangle}(v) + 4 \cdot \text{triangle}(v) \right) \\
 &= 1 - 4 \cdot \text{edge}(v) + \frac{4}{3} \cdot \text{triangle}(v) + 4 \cdot \text{triangle}(v)
 \end{aligned}$$

$$\frac{1}{3} \cdot \text{triangle}(v) = \frac{1}{n} \sum_{v \in V(G)} \text{triangle}(v)$$



$$\text{triangle}(v) = \frac{1}{n} \sum_{v \in V(G)} \text{triangle}(v)$$

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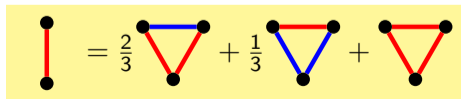
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$$\begin{aligned}
 0 &\leq \frac{1}{n} \sum_v \left(1 - 2 \cdot \text{edge}_v \right)^2 = \frac{1}{n} \sum_v \left(1 - 4 \cdot \text{edge}_v + 4 \cdot \text{triangle}_v + 4 \cdot \text{triangle}_v \right) \\
 &= 1 - 4 \cdot \text{edge} + \frac{4}{3} \cdot \text{triangle} + 4 \cdot \text{triangle}
 \end{aligned}$$





A diagrammatic equation on a yellow background showing an edge (two red lines) equal to the sum of two triangles (one with two red and one blue edge, and one with three red edges) plus another triangle with three red edges.

EXAMPLE - MANTEL'S THEOREM

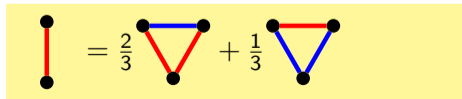
THEOREM (MANTEL 1907)

A triangle-free n -vertex graph contains at most $\frac{1}{4}n^2 \approx \frac{1}{2}\binom{n}{2}$ edges.

Assume edges are red and non-edges are blue.

Assume  = 0. (We want to conclude  $\leq \frac{1}{2}$.)

$$\begin{aligned}
 0 &\leq \frac{1}{n} \sum_v \left(1 - 2 \cdot \text{edge}(v) \right)^2 = \frac{1}{n} \sum_v \left(1 - 4 \cdot \text{edge}(v) + 4 \cdot \text{triangle}(v) + 4 \cdot \text{triangle}(v) \right) \\
 &= 1 - 4 \cdot \text{edge} + \frac{4}{3} \cdot \text{triangle}
 \end{aligned}$$





$$\text{edge} = \frac{2}{3} \cdot \text{triangle}_1 + \frac{1}{3} \cdot \text{triangle}_2$$





EXAMPLE - MANTEL'S THEOREM



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

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$$0 \leq \frac{1}{n} \sum_v \left(1 - 2 \text{}_v \right)^2 = \frac{1}{n} \sum_v \left(1 - 4 \text{}_v + 4 \text{}_v + 4 \text{}_v \right)$$

$$= 1 - 4 \text{} + \frac{4}{3} \text{}$$

$$0 = 2 \text{} - \frac{4}{3} \text{} - \frac{2}{3} \text{}$$



$$\text{} = \frac{2}{3} \text{} + \frac{1}{3} \text{}$$

EXAMPLE - MANTEL'S THEOREM

THEOREM (MANTEL 1907)

A triangle-free n -vertex graph contains at most $\frac{1}{4}n^2 \approx \frac{1}{2}\binom{n}{2}$ edges.

Assume edges are red and non-edges are blue.

Assume  = 0. (We want to conclude  $\leq \frac{1}{2}$.)

$$0 \leq \frac{1}{n} \sum_v \left(1 - 2 \cdot \text{red edge } v \right)^2 = \frac{1}{n} \sum_v \left(1 - 4 \cdot \text{red edge } v + 4 \cdot \text{red triangle } v + 4 \cdot \text{red triangle } v \right)$$

$$= 1 - 4 \cdot \text{red edge} + \frac{4}{3} \cdot \text{red triangle}$$

$$= 1 - 2 \cdot \text{red edge} - \frac{2}{3} \cdot \text{blue triangle}$$

$$0 = 2 \cdot \text{red edge} - \frac{4}{3} \cdot \text{red triangle} - \frac{2}{3} \cdot \text{blue triangle}$$



$$\text{red edge} = \frac{2}{3} \cdot \text{red triangle} + \frac{1}{3} \cdot \text{blue triangle}$$

EXAMPLE - MANTEL'S THEOREM

THEOREM (MANTEL 1907)

A triangle-free n -vertex graph contains at most $\frac{1}{4}n^2 \approx \frac{1}{2}\binom{n}{2}$ edges.

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$$0 \leq \frac{1}{n} \sum_v \left(1 - 2 \cdot \text{edge}_v \right)^2 = \frac{1}{n} \sum_v \left(1 - 4 \cdot \text{edge}_v + 4 \cdot \text{triangle}_v + 4 \cdot \text{triangle}_v \right)$$

$$= 1 - 4 \cdot \text{edge}_v + \frac{4}{3} \cdot \text{triangle}_v$$

$$= 1 - 2 \cdot \text{edge}_v - \frac{2}{3} \cdot \text{triangle}_v$$

$$\leq 1 - 2 \cdot \text{edge}_v$$

$$0 = 2 \cdot \text{edge}_v - \frac{4}{3} \cdot \text{triangle}_v - \frac{2}{3} \cdot \text{triangle}_v$$

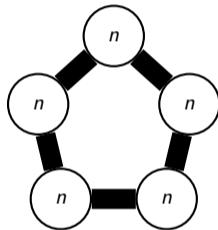
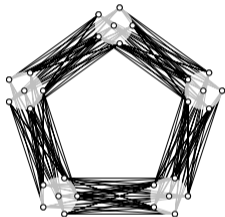
$$\text{edge}_v = \frac{2}{3} \cdot \text{triangle}_v + \frac{1}{3} \cdot \text{triangle}_v$$

Erdős Pentagon Problem

PENTAGONS IN TRIANGLE-FREE GRAPHS

PROBLEM (ERDŐS, 83)

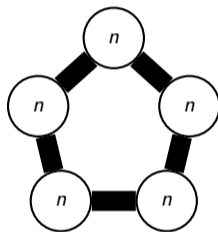
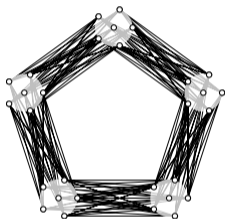
Is it true that a triangle-free graph on $5n$ vertices can contain at most n^5 pentagons?



PENTAGONS IN TRIANGLE-FREE GRAPHS

PROBLEM (ERDŐS, 83)

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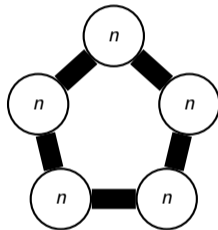
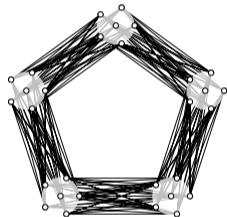
THEOREM (GRZESIK '12 & HATAMI, HLADKÝ, KRÁL', NORIN, RAZBOROV '13)

For all $n > n_0$ or $5|n$, the balanced blow-up of C_5 maximizes the number of C_5 s over all triangle free graphs, and it is unique.

PENTAGONS IN TRIANGLE-FREE GRAPHS

PROBLEM (ERDŐS, 83)

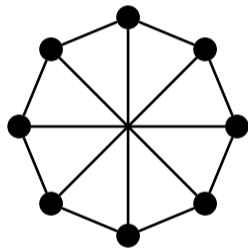
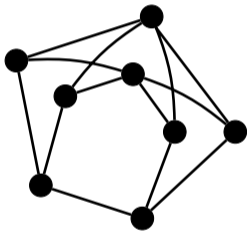
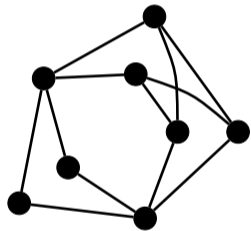
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THEOREM (GRZESIK '12 & HATAMI, HLADKÝ, KRÁL', NORIN, RAZBOROV '13 & L., PFENDER '18)

For all $n \geq n_0$ or $5 \mid n$, the balanced blow-up of C_5 maximizes the number of C_5 s over all triangle free graphs, and it is unique unless $n < 5$ or $n = 8$.

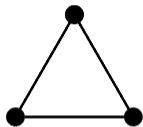
EXTREMAL EXAMPLES ON 8 VERTICES



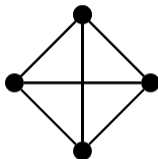
EXTENSIONS OF PENTAGON PROBLEM

PROBLEM (PALMER, 2018)

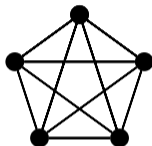
Which K_r -free graph on n vertices contains the most pentagons?



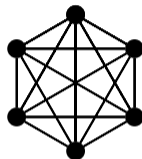
K_3



K_4



K_5

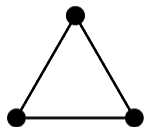


K_6

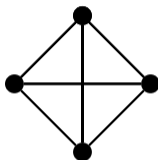
EXTENSIONS OF PENTAGON PROBLEM

PROBLEM (PALMER, 2018)

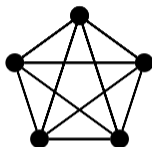
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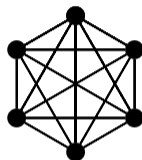
K_3



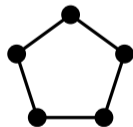
K_4



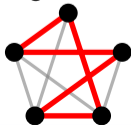
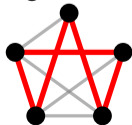
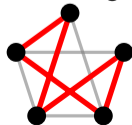
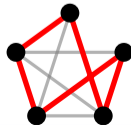
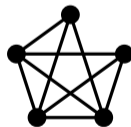
K_5



K_6



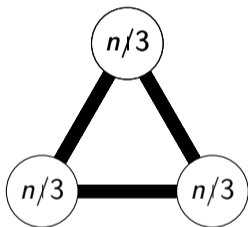
vs



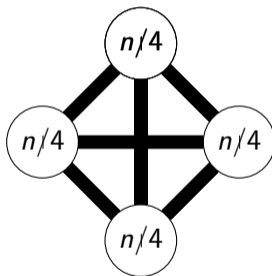
EVERY C_5 COUNTS

THEOREM (L., MURPHY (2021))

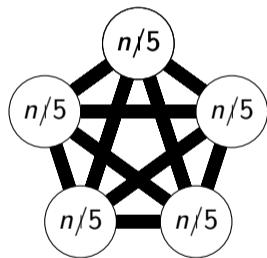
For all $r \geq 3$, the number of 5-cycles among K_{r+1} -free graphs is maximized by the Turán's graph $T_r(n)$ for n sufficiently large.



$T_3(n)$



$T_4(n)$



$T_5(n)$

EVERY C_5 COUNTS

THEOREM (L., MURPHY (2021))

For all $r \geq 3$, the number of 5-cycles among K_{r+1} -free graphs is maximized by the Turán's graph $T_r(n)$ for n sufficiently large.

Flag Algebras formulation:

Maximize

Subject to $K_{r+1} = 0$

MAXIMIZING OTHER GRAPHS IN K_r -FREE

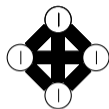
$\text{ex}(n, \#H, F) :=$ Maximum number of copies of H in F -free graph on n vertices.



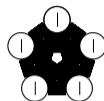
$T_2(n)$



$T_3(n)$



$T_4(n)$



$T_5(n)$

THEOREM (MANTEL (1907))

$\text{ex}(n, \#K_2, K_3) = |E(T_2(n))|$. Moreover, $T_2(n)$ is the unique extremal graph.

THEOREM (TURÁN (1941))

$\text{ex}(n, \#K_2, K_{r+1}) = |E(T_r(n))|$ for $r \geq 3$, and $T_r(n)$ is the unique extremal graph.

THEOREM (ERDŐS-STONE (1946), ERDŐS-SIMONOVITS (1966))

$$\text{ex}(n, \#K_2, F) = \left(1 - \frac{1}{\chi(F)-1}\right) \frac{n^2}{2} + o(n^2).$$

MAXIMIZING OTHER GRAPHS IN K_r -FREE

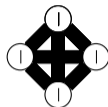
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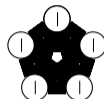
$T_2(n)$



$T_3(n)$



$T_4(n)$



$T_5(n)$

THEOREM (ZYKOV (1949))

Let $t \leq r$. $\text{ex}(n, \#K_t, K_{r+1})$ is maximized in $T_r(n)$.

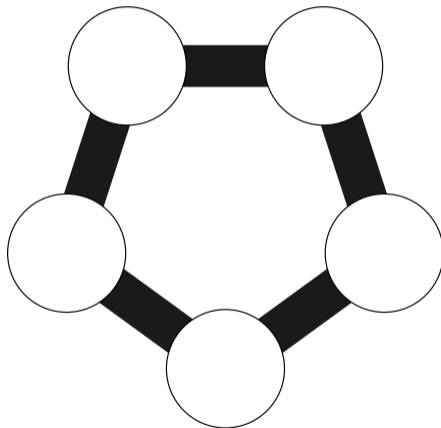
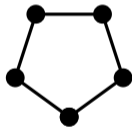
THEOREM (ALON, SHIKHELMAN (2015))

$$\text{ex}(n, \#K_3, C_5) \leq (1 + o(1)) \frac{\sqrt{3}}{2} n^{3/2}$$

Recent results by Gerbner+Palmer, Ma+Qui, Qian+Xie+Ge, Murphy+Nir

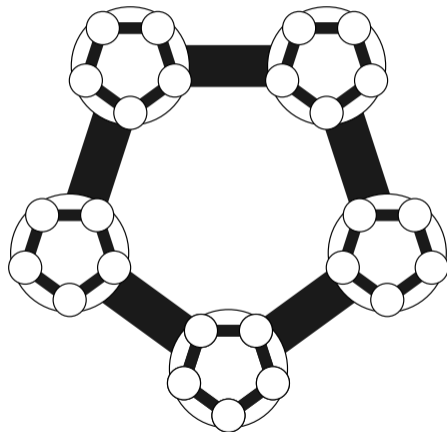
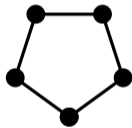
MAXIMIZING PENTAGONS (INDUCED VERSION)

Maximize



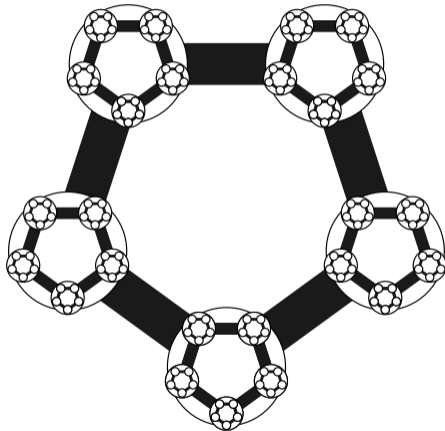
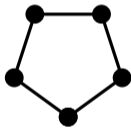
MAXIMIZING PENTAGONS (INDUCED VERSION)

Maximize



MAXIMIZING PENTAGONS (INDUCED VERSION)

Maximize



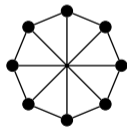
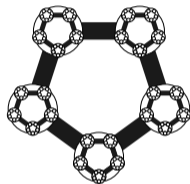
RESULTS

THEOREM (BALOGH, HU, L., PFENDER, 2016)

The iterated blow-up of C_5 maximizes the number of 5-cycles on 5^n vertices.

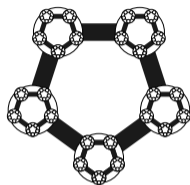
THEOREM (L., MATTES, PFENDER, 2023)

The iterated blow-up of C_5 maximizes the number of 5-cycles on n vertices. Except $n = 8$.



FRACTALIZERS

A graph G is a *fractalizer* if the graph maximizing the number of induced copies of G is an iterated blow-up of G .

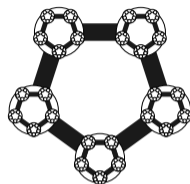


THEOREM (FOX, HUANG, LEE 2015+, YUSTER 2019)

Almost every graph is a fractalizer.

FRACTALIZERS

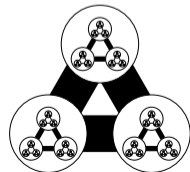
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THEOREM (FOX, HUANG, LEE 2015+, YUSTER 2019)

Almost every graph is a fractalizer.

Can you find some? Other than K_n or $\overline{K_n}$.



CONJECTURE (PIPPINGER, GOLUMBIC 1975)

Cycles, except 4-cycle, are fractalizers.

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Cycles, except 4-cycle, are fractalizers.

THEOREM (L., MATTESS, PFENDER 2023)

5-cycle is almost a fractalizer; exception on 8 vertices.



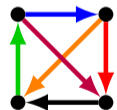
THEOREM (BLUMENTHAL, PHILLIPS, 2021+)

Net is a fractalizer if on 6^k vertices.



THEOREM (MUBAYI, RAZBOROV 2021)

Rainbow tournaments on at least 4 vertices are fractalizers.



Graph fractalizers need at least 8 vertices.

Crossing numbers

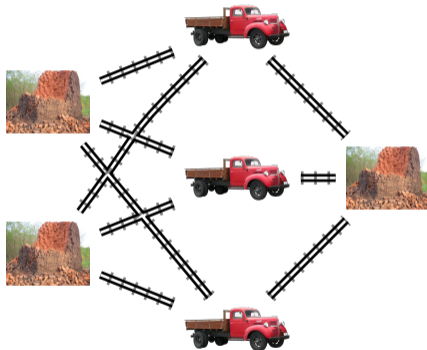
Application to graph drawing.

CROSSING NUMBER

Turán 1945: In a forced labor camp, prisoners transfer carts of bricks from kilns to shipping yards.

When two tracks cross, cart is likely to derail.

How to connect every kiln and shipping yard that minimizes the number of crossings?



$K_{m,n}$ is a complete bipartite graph with sizes m and n , $K_{3,3}$ is above.

CROSSING NUMBER

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$K_{m,n}$ is a complete bipartite graph with sizes m and n , $K_{3,3}$ is above.

For a graph G , $cr(G)$ is the *crossing number*.

CONJECTURE (ZARANKIEWICZ 1954)

$$cr(K_{m,n}) = \left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{(n-1)}{2} \right\rfloor \left\lfloor \frac{m}{2} \right\rfloor \left\lfloor \frac{(m-1)}{2} \right\rfloor$$

THEOREM (NORIN, ZWOLS 2013)

$$cr(K_{m,n}) \geq 0.905 \left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{(n-1)}{2} \right\rfloor \left\lfloor \frac{m}{2} \right\rfloor \left\lfloor \frac{(m-1)}{2} \right\rfloor$$

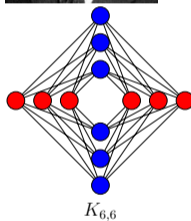
for large m and n . (Zarankiewicz's conjecture is 90.5% true)

80% Kleitman 1970

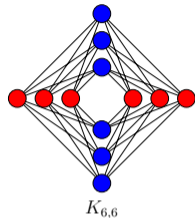
83% De Klerk, Maharry, Pasechnik, Richter, Salazar 2006

85.9% De Klerk, Pasechnik, Schrijver 2007

88.7% Brosch, Polak 2022+



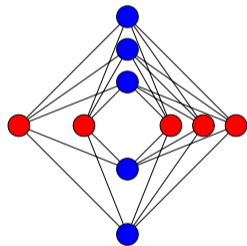
Rectilinear crossing number is with straight line drawing.



THEOREM (BALOGH, L., NORIN, PFENDER, SALAZAR, SPIRO 2023+)
Zarankiewicz's conjecture is 91.1% true for large m and n .

THEOREM (BALOGH, L., NORIN, PFENDER, SALAZAR, SPIRO 2023+)
Rectilinear version of the Zarankiewicz is 97.3% true for large m and n .

$K_{5,5,5}$



$K_{5,5,5}$

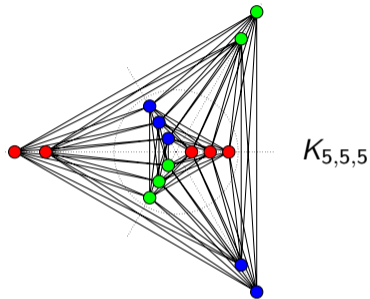
$K_{5,5,5}$

$K_{5,5,5}$

For a graph G , $\overline{cr}(G)$ is the rectilinear crossing number.

CONJECTURE

$\overline{cr}(K_{n_1, n_2, n_3})$ is minimized by

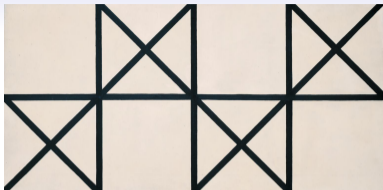


THEOREM (GETHNER, HOGBEN, L., PFENDER, RUIZ, YOUNG, '17)

$\overline{cr}(K_{n_1, n_2, n_3})$ conjecture is 89% true for large n_1 , n_2 , and n_3 .

PROBLEM

What about partite graphs with more parts?

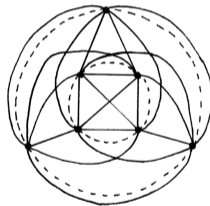


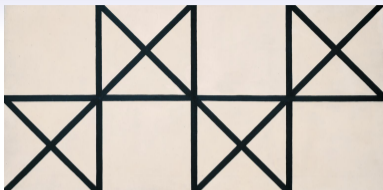
Anthony Hill
Orthogonal / Diagonal
Composition
1954



Hill considered crossing
number of complete graphs.

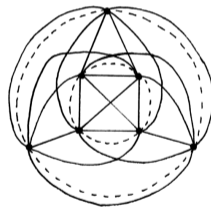
Anthony Hill
Orthogonal / Diagonal
Composition
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Orthogonal / Diagonal
Composition
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Hill considered crossing
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CONJECTURE (HILL 1962)

$$cr(K_n) = \frac{1}{4} \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor \lfloor \frac{n-2}{2} \rfloor \lfloor \frac{n-3}{2} \rfloor$$

CONJECTURE (HILL 1962)

$$cr(K_n) = \frac{1}{4} \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor \lfloor \frac{n-2}{2} \rfloor \lfloor \frac{n-3}{2} \rfloor$$

Conjecture is true

- if $n \leq 12$.
- 100% with various additional restrictions on the drawing
- 80% Kleitman 1970
- 83% De Klerk, Maharry, Pasechnik, Richter, Salazar 2006
- 85.9% De Klerk, Pasechnik, Schrijver 2007
- 90.5% Norin, Zwols 2013
- 91.1% Balogh, L., Norin, Pfender, Salazar, Spiro 2023

CONJECTURE (HILL 1962)

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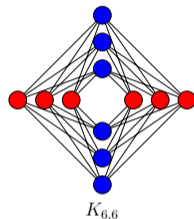
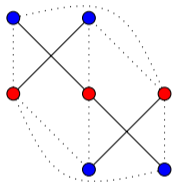
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THEOREM (BALOGH, L., SALAZAR 2019)

Conjecture is 98.5% true.

TUÁN TYPE RESULT



THEOREM (BALOGH, L., NORIN, PFENDER, SALAZAR, SPIRO 2023+)

If $K_{n,n}^D$ is a drawing of $K_{n,n}$ where no $K_{3,4}$ induces exactly two crossings sharing one vertex, then $K_{n,n}^D$ has at least $n^4/16 + o(n^4)$ crossings. (100% True)

HOW IS IT DONE?

For multipartite graphs

- Color vertices to indicate parts.
- For every $\{(a, b), (c, d)\}$, where a, b, c, d are vertices remember if edges ab and cd cross or not.
- Necessary to generate all (combinatorial) embeddings of graphs on n vertices.

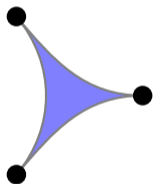
For complete graphs

- For every vertex remember clockwise order of its neighbors.
- Necessary to generate all (combinatorial) embeddings of graphs on n vertices.

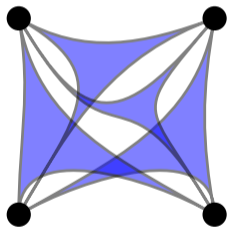
ℓ_2 -norm in Turán Type Problems

HYPERGRAPH SETTING

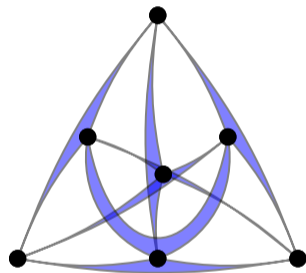
3-uniform hypergraphs have triples of vertices as edges.



One Edge, K_3^3



Complete hypergraph K_4^3



Fano Plane

Turán's Tetrahedron problem: Determine $ex(n, \#K_3^3, K_4^3)$

\$500 reward by Erdős

TURÁN'S TETRAHEDRON PROBLEM

Determine $\text{ex}(n, \#K_3^3, K_4^3)$

Asymptotic setting:

$$\pi(K_4^3) = \lim_{n \rightarrow \infty} \text{ex}(n, \#K_3^3, K_4^3) / \binom{n}{3}$$



K_3^3



K_4^3

THEOREM (KOSTOCHKA 1982, BROWN 1983,
FON-DER-FLAASS 1988, FROHMADE 2008)

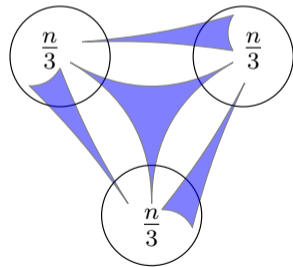
$$\pi(K_4^3) \geq 5/9$$

THEOREM (BABER 2012)

$$\pi(K_4^3) \leq 0.5615$$

THEOREM (RAZBOROV 2010)

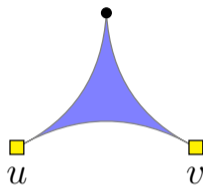
$$\pi(K_4^3, \text{few other graphs}) = 5/9$$



CO-DEGREE VECTOR

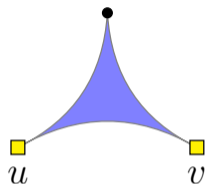
Let G be an n -vertex 3-uniform hypergraph
co-degree vector $X \in \mathbb{Z}^{\binom{n}{2}}$ is indexed by pairs $u, v \in V(G)$
 $X_{u,v} := \#$ edges containing u and v .

$$\# \text{edges} = \frac{1}{3} \sum_{uv} X_{uv} = \frac{1}{3} \|X\|_1$$



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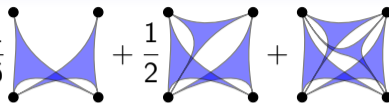
$$\# \text{edges} = \frac{1}{3} \sum_{uv} X_{uv} = \frac{1}{3} \|X\|_1$$

New idea: Consider

$$\sum_{uv} X_{uv}^2 = (\|X\|_2)^2$$

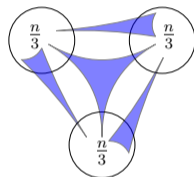
$$\frac{1}{\binom{n}{2}} \sum_{uv} \left(\text{hyperedge}(u,v) \right)^2 = \frac{1}{6} \left(\text{hyperedge}(u,v) \cap \text{hyperedge}(u,v) \right) + \frac{1}{2} \left(\text{hyperedge}(u,v) \cap \text{hyperedge}(u,v) \right) + \left(\text{hyperedge}(u,v) \cap \text{hyperedge}(u,v) \right)$$

Let \mathcal{G}_n be H -free 3-uniform hypergraphs on n vertices.

$$\sigma(H) := \lim_{n \rightarrow \infty} \max_{G \in \mathcal{G}_n} \frac{1}{6} \text{ (diagram 1) } + \frac{1}{2} \text{ (diagram 2) } + \text{ (diagram 3) }$$


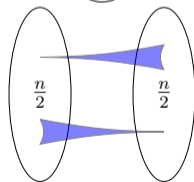
THEOREM (BALOGH, CLEMEN, L. 2022)

$$\sigma(K_4^3) = \frac{1}{3}$$



THEOREM (BALOGH, CLEMEN, L. 2022)

$$\sigma(K_5^3) = \frac{5}{8}$$



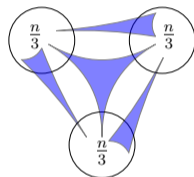
Any many others.

Let \mathcal{G}_n be H -free 3-uniform hypergraphs on n vertices.

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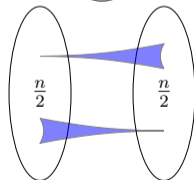
THEOREM (BALOGH, CLEMEN, L. 2022)

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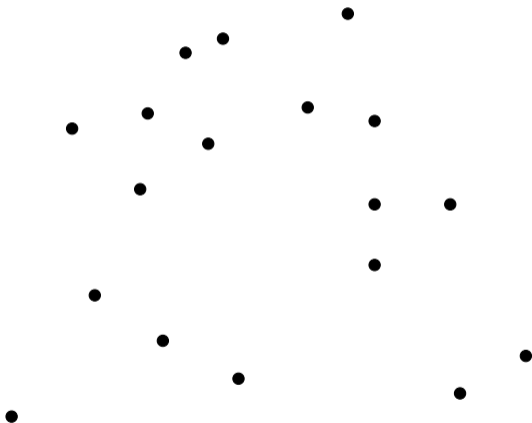
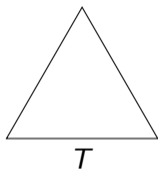
Future: Better exactness methods and other settings applications.

ε -similar Triangles

PROBLEM

Let T be a triangle and $n \in \mathbb{N}$ fixed.

Which n points in \mathbb{R}^2 maximize the number of triangles similar to T ?

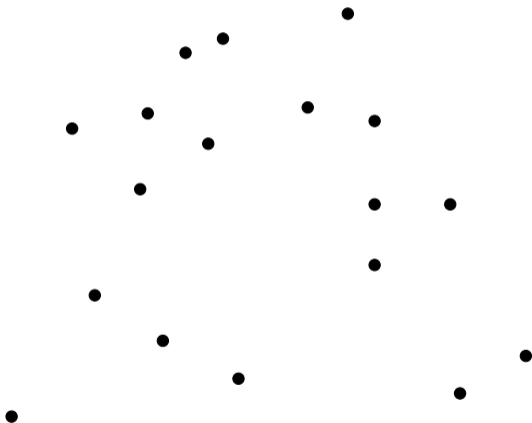
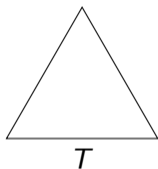


T_1 and T_2 are *ε -similar* if their inner angles differ by at most ε .
(OK to move, scale, rotate, ε -perturb)

PROBLEM

Let T be a triangle and $n \in \mathbb{N}$ fixed. (and $\varepsilon > 0$ fixed)

Which n points in \mathbb{R}^2 maximize the number of triangles similar to T ?

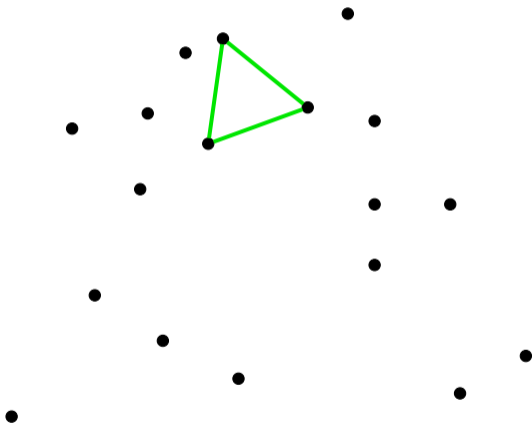
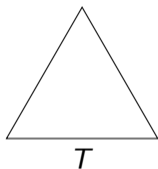


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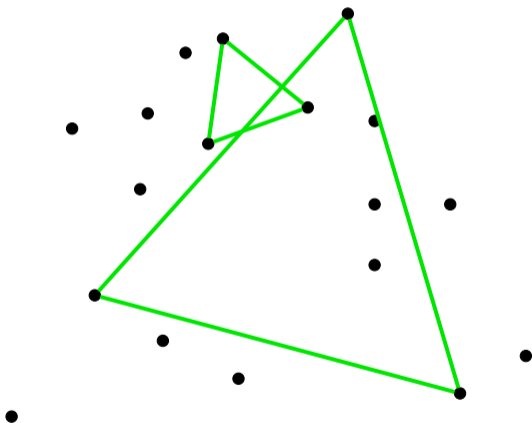
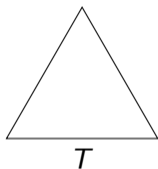


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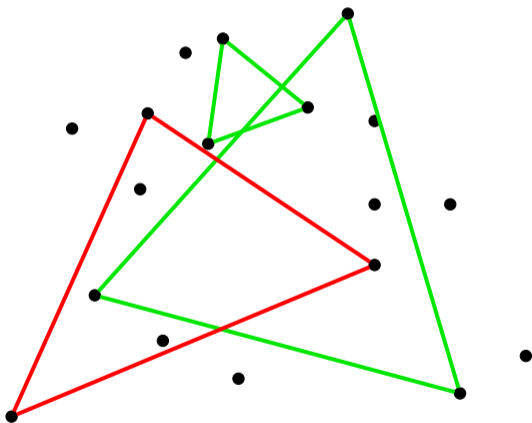
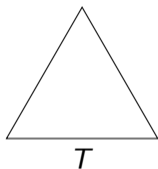


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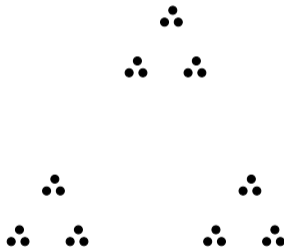
LOWER BOUND CONSTRUCTION

Let T be a triangle and $n \in \mathbb{N}$ fixed. (and $\varepsilon > 0$ fixed)

Which n points in \mathbb{R}^2 maximize the number of triangles similar to T ?



T



$h(n, T, \varepsilon) := \max \#$ of ε -similar triangles to T , it is at least $\frac{1}{4} \binom{n}{3} (1 + o(1))$.

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RESULTS

THEOREM (BÁRÁNY AND FÜREDI (2019))

For almost every triangle T there is an $\varepsilon_0 > 0$ such that for all $0 < \varepsilon \leq \varepsilon_0$,

$$h(n, T, \varepsilon) \leq 0.25072 \binom{n}{3} (1 + o(1)).$$

If T is equilateral, then $h(n, T, \varepsilon) = \frac{1}{4} \binom{n}{3} (1 + o(1))$

THEOREM (BALOGH, CLEMEN, L. (2022))

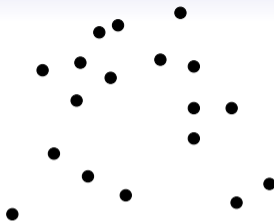
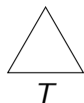
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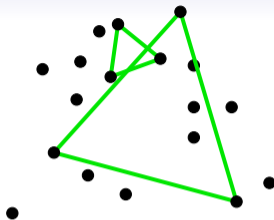
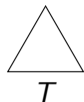
Let T and ε are given

- Fix n points in the plane.
- For every T' ε -similar to T , add a 3-edge
- Investigate the resulting hypergraph H
 H has no subhypergraph in $\mathcal{F} = \{K_4^3, \dots\}$



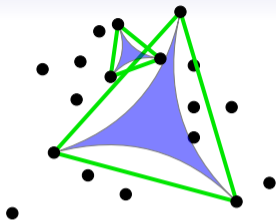
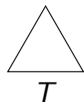
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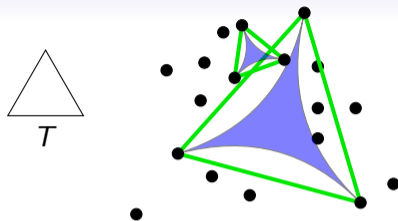


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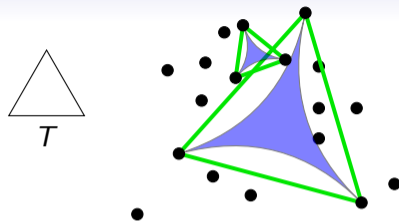
THEOREM (BALOGH, CLEMEN, L. (2022))

\mathcal{F} -free hypergraph has at most $\frac{1}{4} \binom{n}{3} (1 + o(1))$ edges.



Let T and ε are given

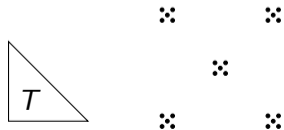
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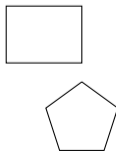
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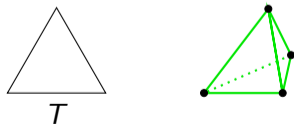
All triangles?



Other Shapes?



in \mathbb{R}^d ?

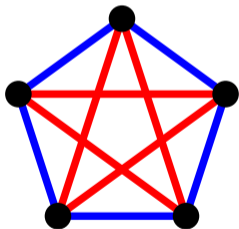


Small Ramsey numbers

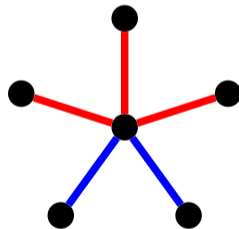
DEFINITION

$R(G_1, G_2, \dots, G_k)$ is the smallest integer n such that any k -edge coloring of K_n contains a copy of G_i in color i for some $1 \leq i \leq k$.

$$R(K_3, K_3) > 5$$



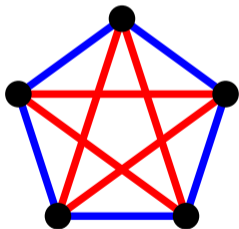
$$R(K_3, K_3) \leq 6$$



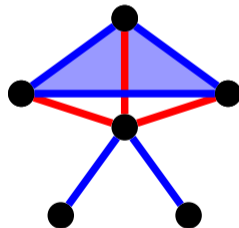
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$$R(K_3, K_3) \leq 6$$



THEOREM (RAMSEY 1930)

$R(K_m, K_n)$ is finite.

$R(G_1, \dots, G_k)$ is finite

Questions:

- study how $R(G_1, \dots, G_k)$ grows if G_1, \dots, G_k grow (large)
- study $R(G_1, \dots, G_k)$ for fixed G_1, \dots, G_k (small)



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Radziszowski - *Small Ramsey Numbers*
Electronic Journal of Combinatorics - Survey



[Erdős] Suppose aliens invade the earth and threaten to obliterate it in a year's time unless human beings can find the Ramsey number for red five and blue five. We could marshal the world's best minds and fastest computers, and within a year we could probably calculate the value. If the aliens demanded the Ramsey number for red six and blue six, however, we would have no choice but to launch a preemptive attack.



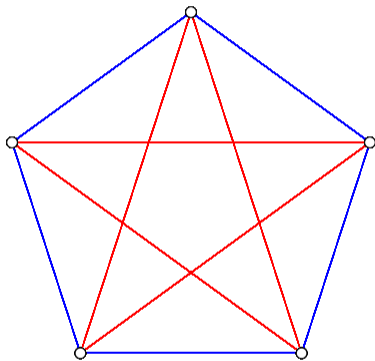
Take any graph G with no





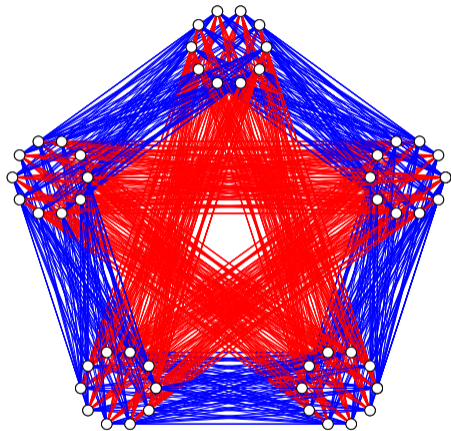
and





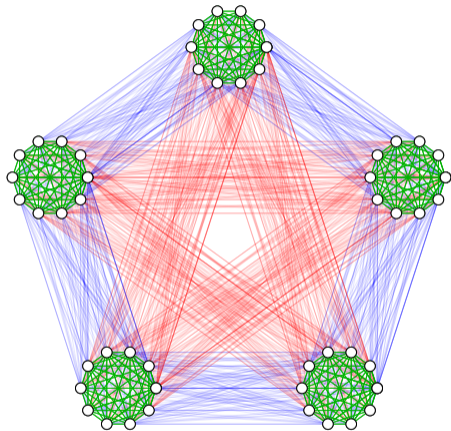
.





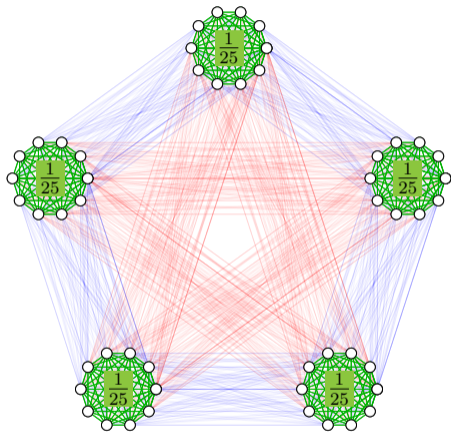
Take any graph G with no  and .





Take any graph G with no  and .



Take any graph G with no  and .



If G has k vertices, then the blow-up has density of non-edges $\geq \frac{1}{k}$.

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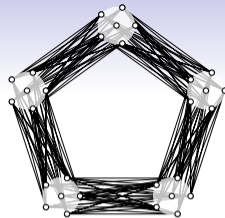
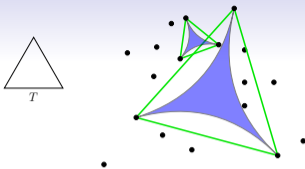
If G has k vertices, then the blow-up has density of non-edges $\geq \frac{1}{k}$.
If any blow-up has density of non-edges $\geq \frac{1}{k}$ then G has $\leq k$ vertices.

NEW UPPER BOUNDS (L., PFENDER 2021)

Problem	Lower	New upper	Old upper
$R(K_4^-, K_8^-)$	29	32	38
$R(K_4^-, K_9^-)$	31	46	53
$R(K_4, K_7^-)$	37	49	52
$R(K_5^-, K_6^-)$	31	38	39
$R(K_5^-, K_7^-)$	40	65	66
$R(K_5, K_6^-)$	43	62	66
$R(K_5, K_7^-)$	58	102	110
$R(K_6^-, K_7^-)$	59	124	135
$R(K_7, K_4^-)$	28	29	30
$R(K_8, K_4^-)$	29	39	42
$R(K_8, C_5)$	29	29	33
$R(K_9, C_5)$	33	36	
$R(K_9, C_6)$	41	41	
$R(K_9, C_7)$	49	58	
$R(K_{2,2,2}, K_{2,2,2})$	30	32	60?

Problem	Lower	New upper	Old upper
$R(K_{3,4}, K_{2,5})$		20	21
$R(K_{3,4}, K_{3,3})$		20	25
$R(K_{3,4}, K_{3,4})$		25	30
$R(K_{3,5}, K_{1,6})$	17	17	
$R(K_{3,5}, K_{2,4})$	16	20	
$R(K_{3,5}, K_{2,5})$	21	23	
$R(K_{3,5}, K_{3,3})$		24	28
$R(K_{3,5}, K_{3,4})$		29	33
$R(K_{3,5}, K_{3,5})$	30	33	38
$R(K_{4,4}, K_{4,4})$	30	49	62
$R(W_7, W_4)$		21	
$R(W_7, W_5)$		16	
$R(W_7, W_6)$		19	
$R(B_4, B_5)$	17	19	20
$R(B_3, B_6)$	17	19	22
$R(B_5, B_6)$	22	24	26

Problem	Lower	New upper	Old upper
$R(W_5, K_6)$	33	36	
$R(W_5, K_7)$	43	50	
$R(Q_3, Q_3)$	13	13	14
$R(K_3, C_5, C_5)$	17	17	21?
$R(K_3, C_4, C_4, C_4)$	24	29	
$R(K_4, C_4, C_4)$	52	71	72
$R(K_4^-, K_4^-, K_4^-)$	28	28	30
$R(K_3, K_4^-, K_4^-)$	21	23	27
$R(K_4, K_4^-, K_4^-)$	33	47	59
$R(K_4, K_4, K_4^-)$	55	104	113
$R(K_3, K_4, K_4^-)$	30	40	41
$R(K_4^-, K_5^-; 3)$	12	12	
$R(K_4^-, K_5; 3)$	14	16	
$R(K_4^-, K_4^-, K_4^-; 3)$	13	14	16



Thank you for your attention!

